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Colin Hinson
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# Interim Technical Document SPB 492 

## TELETEXT SPECIFICATION

## (625 line television systems)

## european broadcasting union

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This document is intended to be the primary reference specification on teletext for European broadcasters. The original information contained was supplied to the EBU by Mr Bernard J. Rogers. Additions have been made to encompass optional codes of practice adopted nationally within certain countries of Europe. As explained in the Introduction, Supplements are available which extend this core documentation to include those features which comprise 'World System Teletext' (also known as CCIR Teletext System B). The EBU acknowledges the considerable contribution to the development of the world System Teletext specification made by Mr Rogers over many years.
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## ABBREVIATIONS

| AL | Address Length |
| :--- | :--- |
| BCD | Binary Coded Decimal |
| CAF | Controlled Access Flag |
| CCIR | International Radio Consultative Committee |
| CEPT | Conference of Postal and Telecommunications Administrations |
| CI | Continuity Indicator |
| CRC | Cyclic Redundancy Check |
| DCLUT | Colour Look-up Table for Dynamically Redefinable Character |
| DIDON | Sets |
| DL | Data Length |
| DRCS | Dynamically Redefinable Character Sets |
| EBU | European Broadcasting Union |
| ECB | Electronic Code Book |
| ESCORT | EBU System of Classification Of Radio and Television |
| PLOF | Frogrammes |
| IAL | Interpretation and Address Length |
| ISO | International Organization for Standardization |
| LUF | Label Up-date Flag |
| LUT | Look-up Table |
| MI | Mode Indicator |
| MJD | Modified Julian Date |
| PDC | Programme Delivery Control |
| PRF | Prepare-to-Record Flag |
| PTU | Pattern Transfer Unit |
| RI | Repeat Indicator |
| TOP | Table of Pages |
| UTC | COordinated Universal Time |

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| [14] - | DIDON uses Layers 1 to 4 of CCIR Teletext system $C$ of [8] |

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## PART A: INTRODUCTION

This specification defines the application of CCIR Teletext System B to CCIR 625 line 50 field Television Systems D, B, G, H, I, K and L [1]. The System is optimised for broadcast media using cable, terrestrial and satellite transmission and the associated service and product environment. Reliable reception of data is ensured, since there is a good match between the service area defined for vision and sound reception and that provided for teletext data broadcasting.

The range of presentation and application features are arranged to be downward compatible. This permits initial services to be defined and existing services to be upgraded, without rendering equipment already in the field obsolete.

The data is organised in a manner optimum for broadcast media by using the rigid timing framework of the television signal. When multiplexed with a video waveform, this permits a fixed relationship to be provided between the data bytes on a television signal data line and locations in the decoder memory. Using this relationship, error identification and correction are available, matched to the statistical occurrence of bit errors.

Critical control data and addressing information are protected by Hamming coding. Basic data is protected by using parity checks. Complete blocks of data may be checked for accuracy using Cyclic Redundancy Check words and these may also be used to provide error correction.

The system provides a presentation layer using alpha-mosaic and DRCS (Dynamically Redefinable Character Sets) character coding. In addition to the characters, a full range of colours can be downloaded providing a colour map with four colour tables, each including eight colours. The presentation features and range of display attributes include those of the CEPT Videotex Presentation Layer Syntax [2], with which they are fully compatible.

A range of display formats is available, from 40 to 160 characters per displayed row and up to 101 rows per page.

The facility to invoke the presentation of characters of any writing system or language, or a mixture of such systems is available. This specification includes the coding for latin based, Greek, Cyrillic and Arabic alphabets. Where appropriate the character repertoires and coding structures of the ISO are used. Associated EBU publications specify the coding for syllabaries [3 \& 4] and idiographic [5] scripts and language requirements that use hybrid combinations of them, plus the coding of musical notes [6].

Alphageometric and Alphaphotographic presentation layers are provided, using the coding methods standardized by the ISO and the service profiles adopted by the CEPT.

In addition to the basic text and graphics display presentation, a wide range of other applications can be supported.

Convenient user access to editorially specified pages and structured groups of pages can be provided by a page linking protocol, transmitted in a special page linking data packet. This permits selected pages to be displayed by a single key stroke, rather than the use of a multi-digit page number. An Inventory Page for each Magazine is provided to optimise memory use in decoders having large storage capacity.

Data relating to the total service on a given channel is carried in a dedicated Broadcast Data Service packet. This packet carries the address of an initial teletext page to be accessed and stored automatically on switch-on or channel selection. Date and time in standardized form are also included with provision for a displayable service message.

To permit applications guch as the automatic control of video recorders, the broadcast service data packet includes data identifying the origin of a broadcast programme, the Network Identification and a programme label. A method of labelling broadcast programmes is provided using a first announced transmission time for compatibility with existing services. The label is intended for association with published programme information using any medium including printed pages, broadcast data and teletext pages. When teletext pages are used, simple programme selection with a cursor is provided by associating machine codes in a dedicated data packet with the displayed items.

The data for transmission can be scrambled to provide access control for Closed User Groups or Subscription User Groups. The descrambling keys can be addressed to individual users or groups of users by an "over air" addressing system, as can credit tokens for "Pay per View" services. Both scrambling and addressing functions are of very high security and reliability.

A parallel EBU publication specifies the application of the system to CCIR 525 line 60 field television systems [7].

## PART B: TRANSMISSION CODING OF TELETEXT IN ANALOGUE COMPOSITECODED TELEVISION SYBTEMS

1. IV Lines Usable as Data Lines

Subject to availability, see Figure 1.
1.1 When Multiplexed with a Composite Video Signal

Lines 7 to 22 and 320 to 335.
Note: In some countries lines 7,318 and 319 may also be used.
1.2 When not Multiplexed with a Composite Video Signal

Any, except field sync and equalising pulse periods.
2. Data Identification

Clock Run-In and Framing Code in appropriate time slot, see Section 9.
3. Signalling Method

Binary NRZ.
4. Data Signal Levels (Negative Modulation)

0 Level Black Level $\pm 2$ \%.
1 Level $66( \pm 6) \%$ of the difference between Black Level and Peak White Level.
5. Bit Rate
$444 \times$ nominal $f_{H}(6.9375 \mathrm{Mbit} / \mathrm{s} \pm 25 \mathrm{ppm})$.
6. Data Timing Reference

Peak Level of penultimate 1 of clock run-in, see Figure 1.
Note: Nominally, the data is inserted with the point of reference at $12.0 \mu \mathrm{~s}$ after the half-amplitude point of the leading edge of the line synchronising pulse. However, it may be necessary to depart from this to allow for the retiming of the synchronising pulses on some Networks, particularly as a result of sync reprocessing in some transmitters, so as to ensure that the timing of the data as transmitted is always within the allowable tolerance range (see also [8]).
7. Spectrum of Data Pulses

Skew symmetrical about $0.5 \times$ bit rate, substantially zero by 5 MHz .
8. Data Line Content

360 bits as 45 bytes of 8 bits each.

## 9. Synchronisation

## See Figure 2.

9.1 Clock Run-In (bit sync)

Bytes 1 and 2, begins 10101010...even parity.

### 9.2 Framing Code (byte sync)

Byte 3, 11100100, even parity.
10. Page Format Data Addressing

See Figures 2 and 3.
10.1 Packet Numbers $X / Y$ When $Y=0$ to $Y=29$

Bytes 4 and 5 Hamming Protected.
3 bits define Magazine Number X.
5 bits define Packet Number Y.
4 bits Hamming Protection in each byte.

### 10.2 Page Header Data Lines

Packet numbers $X / 0$, see Figure 2.

### 10.3 Page Addressing

The Page Address consists of a Page number and a Page Sub-code. The Page Address may take any value as defined in Sections 10.3.1 and 10.3 .2 except:
(i) The address FE 3F7E is reserved for use as the default address for the Magazine Inventory Page, see Section 12.5 .
(ii) The address FF 3F7F is defined as a Null Page Address and shall not be transmitted in a page header packet $Y=0$, see Sections 12.2.2, 12.4.2, 13.2.2, 13.3.2, 15.6, and Appendix 5.

### 10.3.1 Page Number

Bytes 6 and 7 Hamming Protected.
4 bits data plus 4 bits Hamming Protection in each byte.
10.3.2 Page Sub-Code

Bytes 8, 9, 10 and 11 Hamming Protected. 4 bits data plus 4 bits Hamming protection in each byte.
Byte 9 bit 8 is Control Bit C4, see Section 11.1.1, and Figure 2.
Byte 11 bits 6 and 8 are respectively Control Bits C5 and C6, see Sections 11.1.2, 11.1.3, and Figure 2.

Note: See Appendix 5 for Code of Practice for application

### 10.4 Page and Magazine Packet Relationships

Following the Page Header Packet with $Y=0$ of a given page, all subsequent packets with $Y=1$ to $Y=28$ inclusive, from the same magazine relate to that same page.

The transmission of a given page begins with, and includes, it Page Header Packet $(Y=0)$. It is terminated by and excludes the next Page Header Packet ( $Y=0$ ) having the same magazine address but, including the sub-code, a different page address.

Any Packets with $Y=29$ relate to all pages with the included magazine address and not to a specific page.

Any Packets with $Y=30$ and $Y=31$ are not page- or magazinerelated.

## PART C: PRESENTATION LEVEL ONE

11. Basic Level One: 7-Bit Syntax

Decoder Responds to:
a) Packet numbers $\mathrm{X} / 0$ to $\mathrm{X} / 23$.
b) Pages 00 to 99, coded BCD.
c) Page Sub-Codes. The 4 digits can take values in the ranges 0 to 3,0 to 9,0 to 7 and 0 to 9 respectively.
d) Optionally to packet numbers $X / 24$ to $X / 29$ and 8/30.

The use of this 7 -bit syntax is invoked by the data in Packet $\mathrm{X} / 28$, see Figure 11. This data may be omitted for some applications, where the 7-bit syntax is used exclusively or is the default condition.
11.1 Control Bits in Page Header Packet $\mathbf{x / 0}$ (see Figure 2)

C4 to C14 are active on being set to 1 .
Bytes 12 and 13 contain C7 to C14, Hamming protected; for C4 to C6 see Section 10.3.2.
11.1.1 C4 Erase Page

Page to be erased, see Appendix 6.
11.1.2 C5 News Flash

All characters intended for display inset into the TV picture, shall be boxed.

Note: Decoders operating in "mixed" mode also display non-boxed characters.
11.1.3 C6 Sub-Title

All characters intended for display inset into the TV picture, shall be boxed.

Note: Decoders operating in "mixed" mode also display non-boxed characters.
11.1.4 C7 Suppress Header

Header display to be suppressed.
11.1.5 C8 Update Indicator

Following data may be limited to include only updated part of page (see Appendix 6).
11.1.6 C9 Interrupted Sequence

Associated page is not in numerical order of page sequence, allowing the header to be excluded from a rolling header display to avoid discontinuities.
11.1.7 Cl0 Inhibit DisplayData addressed to rows 1 to 24 is not to be displayed.
11.1.8 Cll Magazine SerialMagazines are transmitted one at a time in sequence.
11.1.9 C12, C13, C14
No response, see Section 14.
11.2 Page Display
The conventional text page displays 24 rows of 40 characters,with an optional header row which includes 32 characters. Afurther "commentary row" may be displayed at the top or bottom ofthe page.
11.2.1 Rows Displayable
24, optionally 25 in Data Packet numbers $X / 0$ to $X / 24$, topto bottom of a magazine $X$ page.
11.2.2 Character-Spaces in Rows 1 to ..... 2440; transmitted from left to right.
11.2.3 Character-Spaces in Page Header Row 032; transmitted from left to right.
11.3 Character Bytes
7 bits plus odd parity bit define a display or control characteroccupying a character-space.
11.4 Character Sets for Display
a) $\quad 94$ alphanumeric characters plus SPACE and DELETE as a default GO set, see Figures 15, 21, 22, 23 and 27.
b) 63 block mosaic characters plus SPACE and 32 alphanumeric characters as a default $G 1$ set. The mosaic characters are displayed with their elements either contiguous or separated, see Sections 11.5.4, 11.5.5 and Figure 18.
Selection between character sets is by means of control characters, see Section 11.5 .
11.5 Control Character Set for Spacing Controls Including Display Attributes
Set of 32 control characters, 5 without response at level 1. Thedecoder defaults to specified attributes at the start of eachdisplay row. Some controls have effect immediately, others at thefollowing character-space. The action of a control persists untilthe end of a row or until the transmission of a further controlthat modifies its action. See Figure 20.
11.5.1 Foreground Colour
White; Yellow; Cyan; Green; Magenta; Red; Blue. Invoked with the selection of an alphanumeric or mosaic character set. See Figure 20.
11.5.2 Black Background
Invoked by control character 'Black Background'.
11.5.3 New Background
Control character causes the obtaining foreground colour to be adopted as background colour.
11.5.4 Contiguous Mosaic Graphics
Mosaic blocks adjoin one another.
11.5.5 Separated Mosaic Graphics
Each mosaic block is surrounded by a border of the background colour.
11.5.6 Hold-Mosaic
A Held-Mosaic character is displayed in place of the character SPACE corresponding to a control character. The held-mosaic character is defined only during the period when the mosaic character set is invoked. It is the mosaic character with bit $6=1$ in its code, most recently transmitted before the control character in the 'held- mosaic' position, provided that there has been no intervening change of display mode, normal/double height or alphanumeric/mosaic. The held-mosaic character is displayed in the original contiguous or separated form.
11.5.7 Conceal
Following characters are to be displayed as SPACES until revealed by a decoder or user operation.

### 11.5.8 Flash

Following characters are to be displayed normally or as SPACES in alternation, under the control of a timing device in the decoder.

### 11.5.9 Boxing

Part of a page is to be inserted into the normal television picture. Protection against false operation is provided by double transmission of the control characters, with action taking place between them.
11.5.10 Double Height

Characters are to be stretched vertically to occupy in addition the corresponding character-space in the display row with the next higher address; that row adopts the same display attributes as the previous row and the data in a packet addressed to that row shall be ignored. See also Section 14.5.2.

## 12. Ancillary Text Related Data

### 12.1 Linked Pages Related to a Given Page and Intended for Automatic Storage or Processing in the Decoder <br> Data carried by Packet X/27, see Figure 3. For recommended transmission order of Packets with $Y=26,27$, and 28 see Appendix 1. <br> 12.1.1 Clock Run-In, Framing Code and Packet Address <br> Bytes 1 to 5 as Sections 9 and 10. <br> 12.2 Linked Page Addressing for Exclusively Editorial Applications Designation Codes

Byte 6, 4 bits data plus 4 bits Hamming protection. Codes 0000 to 0011 designate editorially linked page function and are used as sequence labels for a number of Packets $X / 27$, see Figure 3. For an application example see Appendix 4.
12.2.1 Linked Page Addresses, Designation Codes 0000 to 0011

Bytes 7 to 42 as 6 groups of 6 bytes. Each group of 6 bytes defines a linked page address, the groups are numbered 0 to 5 in order of transmission. See also Section 12.2.3.
12.2.2 Data Group Format Defining an Editorially Linked Page

6 bytes containing:
Relative Magazine Number : 3 bits
Page Number: 8 bits
Page Sub-Code : 13 bits
Hamming Protection : 24 bits
For bit sequence see Figure 3.

> When no particular page number is to be specified, the page number FF is transmitted. When no particular page sub-code is to be specified, page sub-code $3 F 7 F$ is transmitted. When the code FF3F7F is transmitted no page is specified. See also Appendix 5 .
12.2.3 Editorial Link Control and Display Row 24 Flag

Byte 43, 4 bits data plus 4 bits Hamming Protection
12.2.3.1 Link Control Data Bit Format

Bit 1 Bit 2

| 0 | 0 | Pages not Chained |
| :--- | :--- | :--- |
| 1 | 0 | Start of Chain of Pages |
| 0 | 1 | End of Chain of Pages |
| 1 | 1 | Within Chain of Pages |
| Bit 3 |  |  |
| 0 |  |  |
| 1 | Data Intended for Direct Display |  |
| Data for Processing of Variable Length |  |  |

For pages of data for processing that are chained, link address 0 with designation code 0000 is the next page in the chain. Other link addresses with designation code 0000 and all link addresses in packets with designation codes 0001 to 0011 specify the addresses of starting pages of chains of pages that include data for processing. Chains may be identified by the designation code and the linked page address.

Data bits 1,2 and 3 above may be reserved for future standardization, when there are no requirements for existing services and products.
12.2.3.2 Display Row 24 Flag

Bit 4:
Set to 'O' Data in packets with $Y=24$ is not to be displayed.

Set to '1' Data in packets with $Y=24$ is to be displayed in row 24.

### 12.3 Basic Page Check Word

Bytes 44 and 45 of packets with the Designation Code 0000 will contain a Cyclic Redundancy Check on data in packets $X / 0$ to $X / 25$.

For Check Word generation see Figure 8.

### 12.4 Linked Page Addressing for Compositional Applications

Byte 6, 4 bits data plus 4 bits Hamming protection. Codes 0100 to 0111 designate linked pages function for compositional applications and are used as sequence labels for a number of packets X/27.

### 12.4.1 Linking Data

Bytes 7 to 42 as 6 groups of $3+3$ bytes, each including 18 bits data plus 6 bits Hamming protection. Bytes 43 to 45 include 18 data bits set to '0' plus 6 bits Hamming protection.
12.4.2 Data Group Format Defining a Compositional Linked Page 6 bytes containing:

| Relative Magazine Number | 3 bits |
| :--- | :--- |
| Page Number | 8 bits |
| Page Sub-Code | 13 bits |
| Link Control Data | 12 bits |
| Hamming Protection | 12 bits |

For bit sequence and interpretation of the Link Control Data, see Figure 3.

When no particular page number is to be specified the page number $F F$ is transmitted. When no particular page sub-code is to be specified page sub-code $3 F 7 F$ is transmitted. When the code FF3F7F is transmitted no page is specified and the link control data bits are set to '1'.

### 12.5 Magazine Inventory Page

This page provides an inventory of the page addresses and number of pages for each address that is currently active in the included magazine. The data is primarily intended to assist memory management in decoders with multipage storage facilities. The data is not intended for direct display, but display can be provided when suitable processing is available.

The data is carried by a page with an address and sub-code defined in an associated packet with $Y=29$ or with the default value FE3F7E (see Figure 6 part 2).
12.5.1 Clock Run-In, Framing Code and Packet Address

Bytes 1 to 5 as Sections 9 and 10.

### 12.5.2 Page Header Packet $Y=0$

### 12.5.2.1 Control Bits in Header Packet $\mathrm{x} / 0$

(see Appendix 6)
Set as follows:

| C4 | Erase Page | Set as required |
| :--- | :--- | :--- |
| C5 | News Flash | Set to 0 |
| C6 | Sub-Title | Set to 0 |
| C7 Suppress Header | Set to 1 |  |
| C8 | Update Indicator | Set as required, |
|  |  | complete page shall |
|  | always be transmitted |  |
| C9 Interrupted Sequence | Set to 1 |  |
| C10 Inhibit Display | Set to 1 |  |
| C11 Magazine Serial | Set as required |  |
| C12, C13, C14 | Have no function in |  |
|  |  | this application and |
|  |  | may have any value. |

### 12.5.2.2 Character-Spaces in Page Header Row 0

Suppress Header Control Bit is to be set, the data bytes corresponding to the 32 character-spaces are reserved.

### 12.5.3 Inventory Data

Carried by packets $X=1$ to $X=16$, bytes 6 to 37 . Bytes 38 to 45 and Packets 17 to 28 are reserved.

```
12.5.3.1 Data Format (see Figure 6 part 1)
    The Data bytes 6 to 37 of packets with Y=1 to
    Y=16 are used as 16 two byte data words. Each
    byte contains }7\mathrm{ active bits plus a parity bit.
    The position of the data word in the packet and
    the packet number define a page address. The
    first data word in the packet with Y=1
    corresponds to page 00. The sequence proceeds
    through the packets with the penultimate data
    word in the packet with Y=16 corresponding to
    page FE.
    The final data word in the packet with Y=16 is a
    continuity indicator, transmitted least
    significant bit first and incremented at every
    change in the content of the inventory page.
```


### 12.5.3.2 Data Word

The 14 data bits, commencing with the least significant bit, are allocated as follows:
(a) Number of pages currently included with the defined page address 13 bits. Maximum number of pages 8190. Set to all $\mathrm{Os}_{8}$, no pages with this address are included in the service. Nevertheless, pages with this address may be transmitted for other applications. Set to hexadecimal 0001, only one version of the page with this address is included, it has sub-code 0000 . Set to hexadecimal 0002, two versions of the page with this address are included, default sub-codes 0001 and 0002. The sequence proceeds to hexadecimal 1FFE, specifying that 8190 versions of the page with this address are included. Set to hexadecimal 1FFF, reserved.
b) Memory Allocation Flag 1 bit

Set to 0, pages include features such as animation and dynamic effects and should not be stored in separate locations. Set to 1, for optimum user access, memory space may be reserved.

## 13. Broadcast Service Data

Packet $8 / 30$, transmitted approximately once per second or more frequently if required by the service.

For bit sequence see Figure 3.

### 13.1 Clock Run-In Framing code and Packet Address

Bytes 1 to 5 as Sections 9 and 10.

### 13.2 Format 1 Packets

Packets shall be transmitted within the field blanking period immediately preceding the boundary between clock-seconds.

### 13.2.1 Designation Code

Byte 6, 4 bits data plus 4 bits Hamming protection:
First data bit set to 0 designates multiplexed function as in Section 1.1. First data bit set to 1 designates non-multiplexed function as in section 1.2. Data bits 2, 3 and 4 set to 0 designate the functions in Sections 13.2.2 to 13.2.8.

| 13.2.2 | Initial Teletext Page for Storage in Decoder without User Action |
| :---: | :---: |
|  | Bytes 7 to 12 containing: |
|  | Absolute Magazine Number 3 bits |
|  | Page Number 8 bita |
|  | Page Sub-Code 13 bits |
|  | Hamming Protection 24 bits |
|  | When no particular page number is to be specified, the page number FF is transmitted. When no particular page sub-code is to be specified the page sub-code $3 F 7 F$ is transmitted. When the code FF3F7F is transmitted, no page is specified. |
| 13.2.3 | Network Identification |
|  | Bytes 13 \& 14: The permanently assigned code uniquely defines the network. |
| 13.2.4 | Time Offset Code |
|  | Byte 15: Defines offset, in half hour units, between local time and Co-ordinated Universal Time (UTC). Negative offsets are West of Greenwich. See Figure 4. |
| 13.2.5 | Modified Julian Date |
|  | Bytes 16 to 18: 5 digit number defining Modified Julian Date (MJD) incrementing daily at midnight UTC. Reference point is 31 January 1982, MJD 45000. See Figure 4. |
| 13.2 .6 | Co-ordinated Universal Time |
|  | Bytes 19 to 21: 6 digit number defining Co-ordinated Universal Time. The transmission relates to the next following second. See Figure 4. |
| 13.2 .7 | lst Short Programme Label |
|  | Bytes 22 and 23: 16 bits define a Programme Label for the currently transmitted programme. |
| 13.2 . 8 | 2nd Short Programme Label |
|  | Bytes 24 and 25: 16 bits define a Programme Label for the currently transmitted programme. |
| 13.2 .9 | Status Display |
|  | Bytes 26 to end of packet: This group is coded with odd parity characters from the default primary character set and where appropriate using the characters common to the range of options. It is intended to display a transmission status message. |

### 13.3 Format 2 Packet

When format 1 packets are also present in a given transmission, the data in bytes 7 to 12 and 26 to the end of the packet shall be the same for both formats. Application information for the data in Section 13.3 is included in [9].

### 13.3.1 Designation Code

Byte 6, 4 bits data plus 4 bits Hamming protection:
First data bit set to 0 designates multiplexed function as in Section 1.1. First data bit set to 1 designates non-multiplexed function as in Section 1.2.
Data bit 2 set to 1 and data bits 3 and 4 set to 0 designate the functions in Sections 13.3.2 to 13.3.7
13.3.2 $\begin{gathered}\text { Initial } \\ \text { Action }\end{gathered}$ Teletext Page for Storage in Decoder without User

Bytes 7 to 12 containing:

| Absolute Magazine Number | 3 bits |
| :--- | ---: |
| Page Number | 8 bits |
| Page Sub-Code | 13 bits |
| Hamming Protection | 24 bits |

When no particular page number is to be specified, the page number FF is transmitted. When no particular page sub-code is to be specified the code 3F7F is transmitted. When the code FF3F7F is transmitted, no page is specified.
13.3.3 Label Channels, Label Up-date Flag and Prepare-to-Record Flag. (see Figure 5 for Bytes 13-26)

Byte 13, 4 bits data plus 4 bits Hamming protection:
Bits 1 and 2 Label Channel Identifier, indicates to which of four parallel data channels the associated label applies.
Data bit 3 Label Update Flag (LUF), set to 'l' indicates that the associated label does not relate to the current television programme, but is intended to update the label memories in video recorders. This provides a method of signalling that the programme which is ending its transmission on a particular channel, is being transferred, perhaps after a period of interruption, to another channel. It may also provide a new label for a programme which is postponed beyond the end of the time window of validity of its original label.
Bit 4 Prepare-to-Record Flag (PRF) indicates to a "waiting" PDC recorder, when set to one, that the programme to which the label applies is about to start but has not yet commenced, and, when reset to zero, that the programme has commenced.
13.3.4. Programme Control Status, Sound Channel Mode and PDC Mode Indicator
Byte 14, 4 bits data plus 4 bits Hamming protection:
bl b2
01 Monophonic Sound
10 Stereo Sound
$1 \quad 1$ Dual Channel SoundBit 3: PDC Mode Indicator (MI), set to 1 indicates thatthe end of transmission of a programme label coincidesexactly with the end of transmission of the programme orthat service codes take immediate effect. When set tozero, it indicates that recording should continue for 30seconds after the programme label is no longer transmitted(and is replaced by another valid label), or that theeffect of service codes is delayed by 30 seconds.
13.3.5 Programme Identification DataBytes 15 to 23, each 4 bits data plus 4 bits Hammingprotection. Comprises data bits 1 to 36.
13.3.5.1 Country Identification Data
Data bits 1 to 4 define the row in the codetable of Figure 29. For column definition seeSection 13.3.5.4.
13.3.5.2 Network Identification or Programme Provider DataData bits 5 and 6 define the final 2 bits of theNetwork Identification or Programme ProviderData Word. For first 6 bits see Section13.3.5.5.
13.3.5.3 Programme Identification Label Data bits 7 to 26 identify a programme by its announced transmission date and time:

| Bits 7 to 11 | Date Day |
| :--- | ---: | :--- |
| Bits 12 to 15 | Date Month |
| Bits 16 to 20 | Time Hours |
| Bits 21 to 26 | Time Minutes |

The complete Label shall be ignored if a value of 3 for the hours units is exceeded when the hours tens has a value of 2 , the value of the hours units exceeds 9, the value of the hours tens exceeds 2. This restriction shall not apply when the month has the value 13 or 14 in which case the complete label consists of an arbitrary number. When the month has the value 15 special application-dependent meaning applies to the complete label.

### 13.3.5.4 Country Identification Data

Data bits 27 to 30 define the column in the code table of Figure 29. For row definition see Section 13.3.5.1.

### 13.3.5.5 Network Provider or Programme Data

Data bits 31 to 36 define the first 6 bits of the national Network Identification or Programme Provider Data Word. The allocation of codes is the responsibility of national authorities. For final 2 bits see Section 13.3.5.2.

```
13.3.6 Programme Type/Series Code Data
Bytes 24 and 25, each 4 bits data plus 4 bits Hamming protection.
The groups of 4 bits represent respectively the first and second digits of a hexadecimal number, each group transmitted least significant bit first. The hexadecimal numbers represent entries in the table of Figure 30 . When all 8 data bits are set to 0 no Programme Type or Series Code is defined.
```


### 13.3.7 Status Display

Bytes 26 to end of packet. See Section 13.2.9.

## PART D: PRESENTATION LEVEL TWO

14. Additional Character Repertoires and Display Attributes

Decoder responds as Level 1 plus Packets $X / 26$ and $X / 28$.
For recommended transmission order of Packets with $Y=26,27$ and 28 (see Appendix 1).
14.1 Control Bits in Page Header Packet $\mathbf{x / 0}$ (see Appendix 6)
14.1.1 C\& TO C11

As Level 1, see Sections 11.1.1 to 11.1.8.
14.1.2 C12, C13 and C14 Primary Character Set Options

Decoder displays text using one of eight options related to the designated or default Primary Character Set, see Figures 14, 17, 25 and 26:

Option Number C12 C13 C14

| $1)$ | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- |
| $2)$ | 0 | 0 | 1 |
| $3)$ | 0 | 1 | 0 |
| $4)$ | 0 | 1 | 1 |
| $5)$ | 1 | 0 | 0 |
| $6)$ | 1 | 0 | 1 |
| $7)$ | 1 | 1 | 0 |
| $8)$ | 1 | 1 | 1 |

14.2 Page Display
14.2.1 Rows Displayed

As Level 1 Section 11.2.1.
14.2.2 Character-Spaces in Display Rows 1 to 24

As Level 1 Section 11.2.2.
14.2.3 Character-Spaces in Display Row 0 Page Header

As Level 1 Section 11.2.3.

### 14.3 Data Bytes

Carried by Packets with $Y=0$ to $Y=29$
14.3.1 Data Packets Other Than Where $Y=26,28$ or 29. As Level 1 Sections 11.3, 12 and 13.
14.3.2 Data Packets Where Y=26, 28 or 29
For these packets, following the clock run-in, framing code, magazine and packet address, there is a designation code in byte 6. The following data is used as a sequence of three-byte groups comprising 18 data bits and 6 Hamming protection bits. For applications see Section 14.6 to end of Section 14.
Packets with $Y=26$ address a character location within a page. They can invoke a character set and write it in the display. They can also invoke a character set and a latching shift. The character sets may be either display or attribute control characters.
Packets with $Y=28$ define the coding method and error protection. They also designate character sets and dynamically redefine the colour map.
Packets with $Y=29$ designate aspects of the display that apply to the magazine defined in the address of the packet with $Y=29$.

### 14.4 Character Sets for Display

Four Character Sets, in GO, G1, G2 and G3 code tables may be designated using the data in packets with $Y=28$, or may be defined as default sets. They are invoked by the data in packets with $Y=26$.
Each designated GO set may have up to 8 options defined. These are invoked by the Page Header Control Bits C12, C13 and C14.
See Figure 14 for character set identification.

### 14.5 Control Character Set for Spacing Controls Including Display Attributes

Set of 32 Control Characters, action as Level 1 , see Section 11.5 , except as defined in Section 14.5, see also Figure 20.
14.5.1 Foreground Colours
As Level 1 plus black, see Section 11.5 .1 and Figure 20.
14.5.2 Size Attributes
Double-Height characters extend downwards, the origin of a character is the upper character position. The double-height and double-size controls are inactive on the bottom row of the defined display area, and the bottom row of a scrolling region.
The whole of a enlarged character is displayed with the attributes that apply to the origin of the character.
Parts of enlarged characters are not displayed: the double-width and double-size controls are inactive in the last character position of a display row.
Attributes set at obscured character positions do not take effect if they would break any of the above rules.

The application of one SIZE attribute control terminates the action of any other SIZE attribute.

At Level 2, where double height or double size characters are included in a given display row, data may be addressed to the row with the next higher address and shall be processed normally, taking into account the above rules. In the absence of such data, the attributes of the previous row are adopted.
14.5.2.1 Double Width

Characters are to be stretched horizontally, to occupy in addition, the next character-space.
14.5.2.2 Double size

Characters are to be stretched horizontally and vertically, to occupy in addition, the character-spaces as in Sections 11.5.10 and 14.5.2.1.

```
14.6 Character Set Extension and Non-Spacing Control Characters for
    Display Attributes
    Uses Packets with Y=26 to overwrite any character-space. The
    original character and attribute condition is the editor-defined
    fallback for Level 1 decoders.
    For recommended transmission order of packets with Y=26, Y=27 and
    Y=28, see Appendix 1.
    14.6.1 Clock Run-In, Framing Code and Packet Address
    Bytes 1 to 5, see Sections 9 and 10.
    14.6.2 Designation Code
    Byte 6, 4 bits data plus 4 bits Hamming protection.
    Codes 0000 to 1111 are sequence labels for up to 16
    packets with Y=26 associated with a given page.
    14.6.3 Data Groups
    Bytes 7 to 45, as 13 groups of 3 bytes each.
```

```
14.6.3.1 Data Group Bit Allocation Format A (Note: other
```

14.6.3.1 Data Group Bit Allocation Format A (Note: other
formats are used for non-alphabetic writing
formats are used for non-alphabetic writing
systems)
systems)
(see Figure 3 part 1)
(see Figure 3 part 1)
6 bits for display address
6 bits for display address
5 bits for mode description
5 bits for mode description
7 bits data
7 bits data
6 bits Hamming protection

```
        6 \text { bits Hamming protection}
```


### 14.6.4 Display Addressing - Display Rows 1 to 24

6 display address bits, as in Section 14.6 .3 provide 64 combinations. The decimal values 0 to 39 specify character-spaces along a given display row. The decimal values 41 to 63 specify a display row, 1 to 23 , value 40 specifies row 24.

A character-space is thus defined explicitly by a data group including a row address, followed by one or more character position data groups.

Although the character-spaces are explicitly addressed, the data groups in Packets with $Y=26$ shall be transmitted in display order, left to right, top to bottom of the display area.

### 14.6.4.1 Display Addressing - Display Row 0 <br> The Row Group Mode Description Bits are set to 00111. The row address is set to correspond to row 23 (decimal value 63). Following data groups with display addresses in the range decimal 8 to 39 specify character-spaces in row 0 .

When the 7 data bits are set to "O" no Full Row Colour is invoked. When data bits 6 and 7 are set respectively to "01", data bits 1 to 5 define the Full Row Colour according to Figure 12. Other settings of bits 6 and 7 and the associated interpretation of bits 1 to 5 are reserved.

The decimal values 40 to 62 shall not be used. Should they occur the associated three byte data group shall be ignored.

Although the character-spaces are explicitly addressed, the data groups in Packets with $Y=26$ of this type shall be transmitted in display order, left to right, of the display area.

### 14.6.5 Display Colours

The Foreground and Background colours of a character cell take precedence over the Full Row and Full Screen colours. Transparent Foreground and Transparent Background colours cause the Full Row or Full Screen colour to be displayed. Full Row colours take precedence over Full Screen colours. Transparent Full Row and Transparent Full Screen permit the display of the video picture where the Foreground and Background colours are also transparent.

### 14.6.5.1 Foreground Colour

For a character displayed at a character space addressed as in Section 14.6.4. The Mode description bits (see Figure 10) are set to 00000. Data bits 1 to 5 define the Foreground Colour according to Figure 12, when data bits 6 and 7 are set to 0. The effect of this attribute persists to the end of a display row unless overridden by a further control defining the Foreground colour.

### 14.6.5.2 Background Colour <br> For a character displayed at a character space addressed as in Section 14.6.4. The mode description bits (see Figure 10) are set to 00011. Data bits 1 to 5 define the Background Colour according to Figure 12, when data bits 6 and 7 are set to 0. The effect of this attribute persists to the end of a display row unless overridden by a further control defining the Background colour.

14.6.5.3 Full Row Colour Including Borders Outside Normal
Text Display Area

Invoked when the mode description bits (see Figure 9) of a row address are set to 00001. Data bits 1 to 5 define Full Row colour according to Figure 12. When data bits 6 and 7 are set to 0, the Full Row colour applies only to the addressed row. When data bits 6 and 7 are set to 1 , the Full Row colour applies to the area from and inclusive of the addressed row to the bottom of the screen, unless overridden by $a$ further Full Row colour control.

### 14.6.5.4 Full Screen Colour Including Borders Outside Normal Text Display Area

Invoked when the mode description bits (see Figure 9) of any row address data group are set to 00000. Data bits 1 to 5 define the Full Screen colour according to Figure 12, when data bits 6 and 7 are set to 0 .

### 14.6.6 Character Sets

Default sets may be specified. Otherwise initial designation is by Packet 28 , see Section 14.8, or Packet 29, see Section 14.9. See also Figures 11 and 14.

Designation is modifiable by the data in Packet 26, see Section 14.6.6.7.

Equipment intended for operation with only a single group of character sets may ignore this data.

```
14.6.6.1 Characters Including Diacritical Marks Composed
    from the Primary and Supplementary Sets
    For display at a character-space addressed as in
    Section 14.6.4. The Mode Description bits, set
    at the range of values }10000\mathrm{ to 11111
    respectively define diacritical marks from
    column 4 of the G2 supplementary character set,
    in ascending numerical order. The associated
    character from the GO primary character set is
    defined by the }7\mathrm{ data bits.
    Character sets are listed in Figure 14.
```

14.6.6.2 Characters from the Supplementary SetFor display at a character-space addressed as inSection 14.6.4. The Mode Description bits areset to 01111. The 7 data bits define a characterfrom the supplementary character set.Character Sets are listed in Figure 14.
14.6.6.3 Block Mosaic Characters (not smoothed)
For display at a character-space addressed as in Section 14.6.4. Mode Description bits are set to 00001. The 7 data bits define a character from the block mosaic character set, when it is the default Gl set or when designated as the Gl set.
14.6.6.4 Block Mosaic Characters (smoothed)For display at a character-space addressed as inSection 14.6.4. Mode Description bits are set to00010. The 7 data bits define a character fromthe smoothed block mosaic character set, when itis the default $G 3$ set or when designated as theG3 set.
14.6.6.5 Latching Shifts to Designated Character SetsFrom a display character-space addressed as inSection 14.6.4, Mode Description bits set to00100 (see Figure 10), the first sixcombinations of the 7 data bits define one ofthe designated character sets. The latchingshift persiats until explicitly cancelled, thetransmission of another group with ModeDescription bits set to 00100 or the end of adisplay row. For other data bit combinations seeSection 14.6.6.6.
14.6.6.6 Combined Character Set Invocation and LatchingFrom a display character-space addressed as inSection 14.6.4, Mode Description bits set to00100 (see Figure 10). Data bit combinations0000110 to 1111110 invoke and latch a characterset defined in Figure 14. The latching shiftpersists until explicitly cancelled, thetransmission of another group with ModeDescription bits set to 00100 or the end of adisplay row.
14.6.6.7 single Shift to Designated Character sets
At a display character-space addressed as in Section 14.6.4, Mode description bits set to 00101 (see Figure 10). The first six combinations of the 7 data bits define one of the designated character sets. For other data bit combinations see Section 14.6.6.8.

### 14.6.6.8 Combined Character Set Invocation and Single shift

At a display character-space addressed as in Section 14.6.4, Mode Description bits set to 00101 see Figure 10. Data bit combinations 0000110 to 1111110 invoke and shift to a character set defined in Figure 14.

### 14.6.6.9 Modified Character Set Designation

The character set designated into the GO, G1, G2 and G3 code tables may be modified from that of packets 28 or 29 by transmitting the Mode Description bits set respectively to 01000 , 01001, 01010 and 01011 see Figure 10). The 7 data bits define a character set from Figure 14. The modified designation persists until the end of a display row or until further modification is signalled. It then reverts to the previously obtaining invocations defined in a packet with $Y=28$ or $Y=29$, or in the absence of such a packet with $Y=28$ or $Y=29$, to the default condition.

### 14.6.7 Display Attributes

These attributes can act together at a given character space. The action persists to the end of a display row and may be cancelled by the transmission of the same data group with the appropriate bit set to '0', or by a level 1 attribute control code. Invoked by the Mode Description bits set to 01100 .

### 14.6.7.1 Mosaic Graphics Separated or Alphanumerics Underlined

Alphanumeric characters following this control character are displayed underlined and mosaic characters are displayed in the separated mode as in Section 11.5 .5 until the receipt of a Cancel-Underline/Contiguous Mosaic Graphics control character or the end of a display row. Activated by data bit 6 set to 1 .

### 14.6.7.2 Boxing/Window

The Boxing attribute acts as Level 1 when the control bits C5 or C6 in the page header packet are set to 1. When neither of these control bits is set to 1 , at Level 2 this attribute has the "window" function. In this case it defines an area where the full screen colour becomes transparent, permitting any video picture to become visible, where the foreground and background colours are also transparent. Activated by data bit 2 set to 1.

### 14.6.7.3 Conceal

Action as Section 11.5 .7 at a character space addressed as in Section 14.6.4. Activated by data bit 3 set to 1 .

### 14.6.7.4 Marked Area <br> The Marked Area starts at a character space addressed as in Section 14.6.4. Activated by data bit 4 set to 1 .

14.6.7.5 Invert

Action at a character space addressed as in Section 14.6.4. This attribute exchanges Foreground and Background Colours and inverts the phase of the flashing clock. Activated by data bit 5 set to 1 .
14.6.7.6 Double Height

Action as Section 11.5.10 at a character-space addressed as in Section 14.6.4. Activated by data bit 1 set to 1 .
14.6.7.7 Double Width

Action as Section 14.5.2.1 at a character-space addressed as in Section 14.6.4. Activated by data bit 7 set to 1 .
14.6.7.8 Double Size

Action as Section 14.5.2.2 at a character-space addressed as in Section 14.6.4. Activated by data bits 1 and 7 both set to 1.
14.6.8 Additional Flash Functions

Action at a character-space addressed as in Section 14.6.4 invoked by mode description bits set to 00111 .

Action of data bits:


The Invert attribute (see Section 14.6.7.5) applies whatever the condition and thus restores an Invert Flash, invoked as above, to Normal, and vice versa.
14.6.9 Scrolling

Invoked when mode description bits of a row address are set to 00101 (See Figure 9).

Data bits 1 to 5 define the colour of the scrolling region from the colour map. Data bits 6 and 7 are set to 0 .

The last row of the scrolling region is defined by the mode description bits of a row address being set to 00110. In this case the data bits 1 to 5 define the Full Row colour from the next row to the border area inclusive. Data bits 6 and 7 are set to 0 .

The scrolling function can only be activated when the page also includes a packet with $Y=27$, with the link control data indicating a chain of pages, see Figure 3 part 3. Data on the first page of a chain for display within the scrolling region and all subsequent pages of the chain may be scrolled through the region under user control. The first page may include data for display above and below the scrolling region. This data is to be displayed throughout the scrolling operation. Data to be scrolled may alternatively be carried on a pseudo page, see Section 15.5. The border of the scrolling region must not be crossed by a double height or a double size character. When the origin of such a character is scrolled out of the scrolling region, the complete character shall disappear from the display.

### 14.6.10 Cursor Display

The display of a cursor is activated by the transmission of the appropriate mode description row group with the bits set to 00100. The column position of the cursor 0 to 39, left to right, is defined by setting the data bits to correspond with the codes $4 / 0$ to $6 / 7$. Other codes are to be interpreted as 'No Cursor'.

### 14.6.11 Termination Marker

Since more than one packet with $Y=26$ may be needed to display a given page, a terminator is provided by setting the Row Address and Mode Description bits all to 1 , in a data group occupying bytes 40,41 and 42 in the final packet with $Y=26$. There is no response to the data in byte 42. Any unused data groups between the active data groups and the termination group shall be filled with repetitions of the data in the termination group.

### 14.6.12 Check Word for Packets with $Y=26$ and $Y=28$

The 18 data bits of the final three-byte data group in the packet with $Y=26$ with designation code 1111, have the two most significant bits set to ' $0^{\prime}$ and a 16 bit Cyclic Redundancy Check on the data in packets with $Y=26$ and $\mathrm{Y}=28$.

Generation of the check word is identical to that of Section 12.3 using the data in packets with $Y=28$ followed by that in packets with $Y=26$. The sequence assumes the presence of 16 packets with $Y=28$ and 16 with $Y=26$. It is completed by assuming that packets that are not present have the 18 data bits in each 3-byte data group set to ' 0 '. When there is data in a packet with $Y=28$ but no data in packets with $Y=26$, only the Termination Marker, (see Section 14.6.11) and the Cyclic Redundancy Check Word will be carried by a packet with $Y=26$.

### 14.7 Colour Dynamic Redefinition

The Colour Map may be redefined, Colour Look-Up Table entries specified and an alternative Colour Table invoked for a given page, using data in packets with $Y=28$. For recommended transmission order of packets with $Y=26,27$ and 28 , see Appendix 1.
14.7.1 Clock Run-In, Framing Code and Packet Address

Bytes 1 to 5 as in Sections 9 and 10.

### 14.7.2 Designation Code

Byte 6, 4 bits data plus 4 bits Hamming protection. Data bits set to 0000 .

### 14.7.3 Data Groups

Bytes 7 to 45 used as 13 groups of 18 bits data and 6 bits Hamming protection.

Bits 1 to 8 of the first group of 18 data bits in these packets with $Y=28$ are set to ' 0 '. Bits 9 to 18 are set according to Figure 11.

The subsequent 12 groups of 18 data bits provide 16 data words of 12 bits each (see Section 14.7.3.1), followed by 4 data words of 5 bits (see Section 14.7.3.2), and a data word of 2 bits (see Section 14.7.3.3), plus two final bits set to '0'.
14.7.3.1 Colour Map Entry Coding

The 16 data words of 12 bits, each define a colour in the Colour Map of Figure 12, proceeding in transmission order from entry 16 to entry 31.

Each 12 bit data word includes 4 bits for each primary colour, red, green and blue, in the transmission order: RRRRGGGGBBBB, with ascending order of bit significance within each 4 bits.
14.7.3.2 Colour Look-Up Table (DCLUT) for Dynamically
Redefinable Character Sets
The 4 data words of 5 bits each define, in order
the 4 entries in the Colour Look-Up Table (DCLUT
of Figure 12), for Dynamically Redefinable
Character Sets (DRCS see Section 16). The
transmission order is least significant bit
first.
14.7.3.3 Invocation of Colour Tables for Use with Spacing Attributes

The 2 bit data word invokes one of the four
Colour Tables to be used with spacing
attributes, see Figure 12. The transmission
order is least significant bit first.
14.8 Designation of Character SetsCharacter Sets are designated to the Code Tables GO, G1, G2 and G3and are defined using the data in packet $\mathrm{x} / 28$.
Equipment intended for operation with only a single group ofcharacter sets may ignore this data.
For the recommended transmission order of packets with $Y=26, Y=27$and $Y=28$, see Appendix 1 .
14.8 .1 Clock Run-In, Framing Code and Packet Address
Bytes 1 to 5 inclusive, see Sections 9 and 10.
14.8.2 Designation Code
Byte 6, 4 bits data plus 4 bits Hamming protection. Databits set to 0001 .
14.8.3 Data Groups
Bytes 7 to 45 used as 13 groups of 18 bits data and 6 bitsHamming protection.
The first two data groups are used for this application.The remaining data groups are reserved.
14.8.3.1 Coding for Character Set Designation - Bit Allocation

        The 18 data bits of the first three-byte data group are allocated as follows:
    (1) Set to ' 0 ' ..... 2 bits GO Table 7 bits, see Figure 14
(3) Set to '0' ..... 1 bit
Character set Code of G1 Table 7 bits, see Figure ..... 14
(5) Set to '0' ..... 1 bit

The 18 data bits of the second three byte data group have the same allocation but (2) designates a G2 Code Table and (4) a G3 Code Table, see Figure 14.

The interpretation of the 7 character code bits is shown in Figure 14.

The bits are transmitted least significant bit first.

In the case of the GO Set, up to 8 options may be designated with the GO set and invoked by the page header control bits C12, C13 and C14, see Section 14.1.2.

The designation of character sets to an associated group of code tables must take into account the relationships between them, avoiding incompatibility.

Some codes in Figure 14 permit simultaneous designation of more than one GO set. Each may have options specified, provided that the total number of options does not exceed 8.
14.9 Data Applicable to All Pages in a Magazine
Data applying to all the pages in a magazine is included in
packets with $Y=29$. Where page applicable data is included in
packets associated with a specifically addressed page, this shall
take precedence over the corresponding data in packets with $Y=29$.
14.9.1 Clock Run-In, Framing Code and Packet Address
Bytes 1 to 5 as Sections 9 and 10.
14.9.2 Designation Code
Byte 6, 4 bits data plus 4 bits Hamming protection.
14.9.3 Colour Dynamic Redefinition
The Designation Code is set to 0000.
The Colour Map may be redefined, Colour Look-Up Table
entries specified and an alternative Colour Table invoked
for all pages in the magazine addressed by a packet with
$\mathrm{Y}=29$. The coding details follow those in Sections 14.7.1
to 14.7.3.3.

### 14.9.4 Designation of Character Sets

    The Designation Code is get to 0001.
    Character sets may be designated to the Code Tables GO,
    G1, G2 and G3 for all pages in the magazine addressed by a
    packet with \(Y=29\). The coding details follow those in
    Sections 14.8.1 to 14.8.3.1.
    
### 14.9.5 Magazine Inventory Page <br> The Designation Code is set to 0100.

### 14.9.5.1 Data Groups

Bytes 7 to 45 as 13 groups of 18 bits data and 6 bits Hamming protection.

### 14.9.5.2 Page Address of Inventory Page

First 2 groups of 3 bytes containing:
Continuity Indicator 2 most significant bits Set to '0' 1 bit Page Number 8 bits Page Sub-Code Continuity Indicator 12 least significant bits Hamming Protection 12 bits

See Figure 6 part 2. The remaining 3-byte data groups are reserved.

The Continuity Indicator shall have the same value as in the associated Inventory Page, see Section 12.5.3.1. When this form of a packet with $Y=29$ is transmitted but no Inventory Page is included, the page and sub-code addresses shall be set to FF3F7F and the Continuity Indicator to 0. The default address for this page when a packet with $Y=29$ is not transmitted is FE3F7E.
15. Pseudo Pages
Pseudo Pages carry data that is intended to be associated with a
standard page or pages. To prevent independent display of a pseudo page
without the associated page, control bits C7 (Suppress Header) and C10
(Suppress Display), in the header packet ( $Y=0$ ) may be set to 1 .
15.1 Linking to Pseudo Pages from an Associated Page
Linking to a pseudo page from the associated page is by means of
the data in packet $X / 27$ of the associated page (see also section
12).
15.2 Clock Run-In, Framing Code and Packet Address
Bytes 1 to 5 as in Sections 9 and 10.
15.3 Overwriting Pseudo Pages
15.3.1 Designation of an Overwriting Pseudo Page
There shall be a packet $\mathrm{X} / 28$ with the Designation Code set
to 0000 .
The first group of 18 data bits shall be set designating
the overwriting mode, see Figure 11. Other data bits set
to '0'.

### 15.3.2 Invocation of Overwriting Modes

The default character set of an Overwriting Pseudo Page is the GO primary character set that is the default, or that designated for standard pages. Other character sets may be invoked for a pseudo page, using packets with $Y=28$, designation code 0001, as in Section 14.8. Character sets are invoked by the transmission of the appropriate Row and Space Mode Description Groups in packets with $Y=26$ of the pseudo page, see Figure 9, 1014 , and Section 14.6. The invoked character set persists until the end of a display row or until overridden by the transmission of a further invocation in a packet with $Y=26$.

### 15.4 Scrolling Data from a Pseudo Page

When the scrolling function of Section 14.6 .9 is activated, the data to be scrolled through the scrolling region may be carried by a pseudo page.
15.4.1 Designation of a Scrolling Pseudo Page

There shall be a packet $\mathrm{X} / 28$ with the Designation Code set to 0000 .

The first group of 18 data bits shall be set designating the scrolling mode, see Figure 11. Other data bits are set to 0 .

The number of pseudo pages required for the scrolling function may be linked by using packet $\mathrm{X} / 27$ of the pseudo page. The associated page introducing the scrolling function is linked as in Section 15.1.

### 15.5 Page Format Extension Using Pseudo Pages

The format of a displayed page may be extended, both horizontally and vertically using data carried by pseudo pages. The horizontal extension can be by up to 3 multiples of 40 characters, i.e. 160 characters per displayed row. The vertical extension can be by up to 3 multiples of 25 rows, i.e the page header plus 100 rows, providing 101 rows per displayed page.
15.5.1 Linking to Extended Format Pages from Standard Pages
Linking to a pseudo page for page extension, from a
standard page is as in Section 15.1 . For this application
the standard page may not include any data for display but
only a suppressed header and one or more packets with
Y=27. It is to be noted that it is not necessary to
introduce an extended page with a standard page. The first
pseudo page of the group forming the extended page can be
acquired directly.

### 15.6 Designation of Pseudo Pages for Page Format Extension

There shall be a packet $\mathrm{X} / 28$ with the Designation Code set to 0000 .

The first group of 18 data bits shall be set designating the page format extension mode, see Figure 11, other data bits are set to 0 . These bits also define the position of the pseudo page in the array forming the complete page.

The termination of the format extension process is defined by the inclusion in the final pseudo page of the group, of a packet with $Y=27$, in which the link control data indicates format extension associated with the NULL page address FF3F7F.

### 15.7 Reformatted Data by Pseudo Pages

An arbitrary data stream may be reformatted into pseudo pages, using the 1024 data bytes available in rows 0 to 25 of a pseudo page. Adequate buffer storage must be available to ensure at least two transmissions of any such pseudo page where the input data stream is subject to dynamic changes.

The reformatting protocols are defined with the specification of the data type involved, or shall be the subject of agreement between the origination and reception ends of the transmission process.

| 15.7.1 | Linking to a Reformatted Data Pseudo Page from a Standard |
| :--- | :--- |
|  | Page |
|  | Linking to a pseudo page carrying data to be reformatted, |
|  | froma standard page is as in Section 15.1 . For this |
|  | application the standard page may not include any data for |
|  | display but only a suppressed header and one or more |
|  | packet with Y=27, see Figure 3 part 3. However it is not |
| necessary to introduce a pseudo page carrying data to be |  |
| reformatted with a standard page. The first pseudo page of |  |
| a linked group may be acquired directly. |  |

There shall be a packet $\mathrm{X} / 28$ with the Designation Code set to 0000 .

The first group of 18 data bits shall be set designating the Reformatted Data mode, see Figure 11. Other data bits set to 0 .

## PART E: PRESENYATION LEVBL THRREE

16. Dynamically Redefinable Character sets - Downloading Using Pseudo Pages
Pseudo Pages of this type include data that define downloadable
character sets for use with a specified page or pages. To prevent
independent display of data intended for downloading, the header
control bits C7 ( Suppress Header) and C1O (Suppress Display), in the
header packet $(Y=0)$ of the pseudo page may be set to 1 .
16.1 Linking from Pages for Display to Pseudo Pages for Character Set Downloading

Linking to a pseudo page for character set downloading is by means
of the data in packet $\mathrm{X} / 27$ of the page for display (see also
Section 12).
16.2 Clock Run-In, Framing Code and Packet Address
Bytes 1 to 5 as in Sections 9 and 10.
16.3 Designation of a Pseudo Page for Character Set Downoading
There shall be a packet with $Y=28$ and the Designation Code set to 0000 .
16.3.1 Mode Definition of Pseudo Pages for Character set Downloading
Two pseudo pages are required for the complete range of DRCS code table addresses.
For code table addresses $2 / 0$ to $4 / 15$ the first group of 18 data bits has bit 8 set to 1 and the remaining bits set to 0 .
For addresses $5 / 0$ to $7 / 15$, bits 8 and 1 are set to 1 and the remaining bits are set to 0 , see Figure 11 part 3.

### 16.4 Character Coding for Downloading

Characters are downloaded using Pattern Transfer Units (PTUs) of 20 bytes each. Each packet X/1 to X/24 carries 2 PTUs in a defined format of 20 plus 20 bytes, proceeding in time through the packet.
16.4.1 Byte Coding for Pattern Transfer Units (PTUs)
The data bytes defining a PTU use the transmission codes 4/0 to 7/15. Each byte thus defines the value of 6 bits of a PTU and is called a D-byte.
16.4.2 $\begin{aligned} & \text { Dynamically Redefinable Character Sets (DRCS) Character } \\ & \text { Modes }\end{aligned}$

| Mode | Format | PTUs/Character | Bytes/Character |
| :--- | :--- | :---: | :--- |
| $(1)$ | $12 \times 10 \times 1$ | 1 | 20 |
| $(2)$ | $12 \times 10 \times 2$ | 2 | 40 |
| $(3)$ | $6 \times 10 \times 1$ | 0.5 | 10 |
| $(4)$ | $6 \times 10 \times 2$ | 1 | 20 |
| $(5)$ | $6 \times 10 \times 4$ | 2 | 40 |
| $(6)$ | $6 \times 5 \times 2$ | 0.5 | 10 |
| $(7)$ | $5 \times 4$ | 1 | 20 |

Format is: Horizontal Dots $x$ Vertical Dots x Bits/Pixel
Up to 96 PTUs may be downloaded using two pseudo pages and these relate directly to the 96 code table addresses for DRCS, 2/0 to 7/15. Unused packets need not be transmitted and unused addresses may be padded with the character SPACE (2/0).

### 16.4.3 Downloading Modes

From the number of PTUs required for the various formats (see Section 16.4.2), the following transmission modes are derived and identified:

| Format |  | Mode Identification <br> 12x10x1 |
| :--- | :--- | :--- |
| $12 \times 10 \times 2$ |  | 0000 |
| $6 \times 5 \times 2$ (1st DRCS) | $6 \times 5 \times 2$ (2nd DRCS) | 0001 |
| $6 \times 10 \times 1$ (1st DRCS) | $6 \times 5 \times 2$ (2nd DRCS) | 0011 |
| $6 \times 5 \times 2$ (1st DRCS) | $6 \times 10 \times 1$ (2nd DRCS) | 0100 |
| $6 \times 10 \times 1$ (1st DRCS) | $6 \times 10 \times 1$ (2nd DRCS) | 0101 |
| $6 \times 10 \times 2$ |  |  |
| $6 \times 10 \times 4$ |  | 0110 |
| $6 \times 5 \times 4$ |  | 0111 |

The Mode Identification Code 1111 indicates that no data is transmitted for the corresponding character and any character already defined persists.

### 16.4.4 Downloading Mode Invocation

The Downloading Mode may be specified individually for each character. The mode for the character is invoked by the second and following groups of 18 data bits in the packet with $Y=28$ associated with the pseudo page or pair of pseudo pages.

The first 192 data bits in the data groups are used to transmit the 48 Mode Identification Codes, see Section 16.4.3, required for each pseudo page used for DRCS downloading. The remaining bits are set to 0 .
16.4.5 DRCS Code Table Organisation

The packet addresses relate directly to positions in the code table.

The firgt PTU in packet X/1 of the first of the pair of downloading pseudo pages, includes dot data for character 2/0. The next PTU includes data for character $2 / 1$ and so on.

The second pseudo page begins with the first PTU of its packet $X / 1$ including dot data for character 5/0 and continuing to character 7/15.

### 16.4.6 Byte Downloading Organisation

The byte organisation in the pseudo page depends upon the character modes and the associated downloading modes to be included in the DRCS, see Figure 13.

| 16.4.6.1 | Character Mode (1) $12 \times 10 \times 1$ (Basic Mode) Mode Identification 0000 <br> The dots of a character are loaded 6 at a time from each D-byte, using the 6 least significant bits from codes $4 / 0$ to $7 / 15$. Loading proceeds from the top left hand corner, left to right, row by row, the most significant bit of each 6 corresponding to the left hand dot. One PTU is required for each $12 \times 10 \times 1$ character. |
| :---: | :---: |
| 16.4.6.2 | Character Mode (2) $12 \times 10 \times 2$ Mode Identification 0001 |
|  | The first bit plane is downloaded as for the basic mode, see Section 14.4.6.1. The second bit plane is downloaded using the next 20 D-byte group. The address in the DRCS code table corresponds to the first 20 D-byte group defining the character. Four colours are available for display, from the DRCS Colour Look-Up Table (DCLUT), downloaded as in Section 14.7.3.2. The first bit plane corresponds to the least significant bit of the DCLUT address and the second bit plane corresponds to the most significant DCLUT address bit. |
| 16.4.6.3 | Character Mode (6) 6x5x2 - Two Co-Defined Character Sets Mode Identification 0010 |
|  | Downloading proceeds as for the basic mode, see Section 16.4.6.1, except that alternate D-bytes respectively define equivalent dots of the first and second sets to be downloaded. Two bit planes are downloaded for each character. Four colours are available for display, from the DRCS Colour Look-Up Table (DCLUT), downloaded as in Section 14.7.3.2. The first bit plane downloaded, corresponds to the least significant bit of the DCLUT address and the second bit plane downloaded corresponds to the most significant DCLUT address bit. |

16.4.6.4 Character Mode (3) $6 \times 10 \times 1$ and Character Mode ..... (6) $6 \times 5 \times 2$ as Two Co-Defined Character Sets ModeIdentification 0011
Downloading proceeds as for the basic mode, see Section 16.4.6.1, except that alternate D-bytes respectively define equivalent dots of the first and second set to be downloaded. For the second of the pair of character sets, two bit planes are downloaded for each character and four colours are available for display, from the DCLUT, downloaded as in Section 14.7.3.2. The first bit plane downloaded corresponds to the least significant bit of the DCLUT address and the second bit plane downloaded corresponds to the most significant DCLUT address bit.
16.4.6.5 Character Mode (6) $6 \times 5 \times 2$ and Character Mode (3) 6x10x1 as Two Co-Defined Character Sets Mode Identification 0100
Downloading is as in Section 16.4.6.4 except that the first and second character set downloading procedure is interchanged.
16.4.6.6 Character Mode (3) $6 \times 10 \times 1$ - Two Co-Defined
Character Sets Mode Identification 0101
Downloading proceeds as for the basic mode, see Section 16.4.6.1, except that alternate D-bytes respectively define equivalent dots of the first and second sets to be downloaded.

### 16.4.6.7 Character Mode (4) $6 \times 10 \times 2$ Mode Identification 0110

Downloading proceeds as for the basic mode, see Section 16.4.6.1, except that the sequence of D-bytes define respectively the least and most significant bits of the DCLUT address. The DCLUT is downloaded as in Section 14.7.3.2.

### 16.4.6.8 Character Mode (5) $6 \times 10 \times 4$ Mode Identification 0111

Two PTUs are required to define a character in this mode. Downloading proceeds as for the basic mode, see Section 16.4.6.1, except that the sequence of D-bytes of the first PTU defines respectively the first and second bit planes. The sequence of D-bytes of the second PTU defines respectively the third and fourth bit planes. The four bit planes define the addresses in colour tables 3 and 4 of the colour map, the first bit plane corresponding to the least significant address bit and so on.

### 16.4.6.9 Character Mode (7) $6 \times 5 \times 4$ Mode Identification 1000

Each group of four D-bytes within the PTU defines respectively the first, second, third and fourth bit planes of a character. The four bit planes define the addresses in colour table 3 and 4 of the colour map, the first bit plane corresponding to the least significant address bit and so on.

## PART F: PRESENTATION LEVEL FOUR

## 17. Alphageometric Displays

There may be an introductory page that is not part of the alphageometric page. The introductory page shall at least include a header packet with $Y=0$ and a packet or packets with $Y=27$, designation codes 0100 to 0111 (see Section 12.4) to provide links to the alphageometric page for display.

These links point to pseudo pages carrying Geometric Instructions and when required pseudo pages for Overwriting. These latter carry alphamosaic characters for insertion into the geometric display. The type of pseudo page is identified by the data in packets with $Y=28$ of those pages.

### 17.1 Pages for Alphageometric Display

The data for display is carried by pseudo pages of two types:
(a) Overwriting Pseudo pages. These pseudo pages are as defined in Section 15.3 and carry alphamosaic character data for association with a geometric display.
(b)

Pages of reformattable data carrying geometric data. Four modes, with increasing display features are defined corresponding to Service Profiles 0, 1, 2 and 3 of [2].

There are two options for the association of the alphamosaic and geometric components of the page:

Option 1 defines the geometric display plane as transparent to the lower alphamosaic character plane.

Option 2 inserts the alphamosaic characters into the plane of the geometric display. The display may thus be considered to be derived from a single plane.

The selection of a mode defining a profile and Option 1 or 2 is by means of the data in the packet with $Y=28$ of the pseudo page carrying the geometric data.

### 17.2 Pseudo Pages Carrying Geometric Data

In a pseudo page carrying geometric data, the header packet with $Y=0$ shall have the Suppress Header control bit $C 7$ and the Inhibit Display control bit C1O both set to 1.

Packets with $Y=1$ to $Y=25$ carry the geometric data according to [2]. The Service Profile, Display option and method of bit coding is defined by the 18 data bits of the first three byte data group of the packet with $Y=28$, see Figure 11.

More than one pseudo page may be required for the geometric data for a displayable page. In this case, the pseudo pages are linked by the data in a packet with $Y=27$.

Unused packets need not be transmitted and incomplete packets shall be filled with the character $0 / 0$.

## PART G: PRESENTATION LEVEL FIVE

## 18. Alphaphotographic Displays

There may be an introductory page that is not part of the alphaphotographic page. The introductory page shall at least include a header packet with $Y=0$ and a packet or packets with $Y=27$, designation codes 0100 to 0111 (see Section 12.4) to provide links to the alphaphotographic page for display.

These links point to pseudo pages carrying Pixel data and, when required, pseudo pages for overwriting. These latter carry alphamosaic characters and geometric patterns for insertion into the photographic display. The type of pseudo page is identified by the data in packets with $Y=28$ of those pages.

### 18.1 Pages for Alphaphotographic Display

The data for display is carried by pseudo pages of three types:
(a) Overwriting Pseudo pages. These pseudo pages are as defined in Section 15.3 and carry alphamosaic character data for association with a photographic display.
(b) Geometric Coded pages for association with a photographic display.
(c) Reformattable pages carrying pixel data. The picture coding method is defined by the data in a packet with $Y=2$ of the pseudo page, according to [10].

There are two options for the association of the alphamosaic, geometric and photographic components of the page:

Option 1 defines the photographic display plane as transparent to the lower alphamosaic and geometric character plane or planes.

Option 2 inserts the alphamosaic characters and geometric pattern into the plane of the photographic display.

The selection of the photographic mode and Option 1 or 2 is by means of the data in the packet with $Y=28$ of the pseudo page carrying the photographic data.

### 18.2 Pseudo Pages Carrying Photographic Data

In a pseudo page carrying photographic data, the header packet with $Y=0$ shall have the Suppress Header control bit $C 7$ and the Inhibit Display control bit clo both set to 1.

Packets with $Y=1$ to $Y=25$ carry the pixel data according to the selected picture coding method.

The Photographic mode, Display Option and method of bit coding is defined by the 18 data bits of the first three byte data group of the packet with $Y=28$, see Figure 11 .
18.2.1 The Method of Picture Coding

The method of picture coding is identified using the second and subsequent three byte data groups in the packet with $Y=28$ of the pseudo page. Two characters according to ISO 9281 are included in each three byte group of 18 data bits, with the final 4 data bits set to ' 0 '.

More than one pseudo page may be required for the pixel data for a displayable page. In this case, the pseudo pages are linked by the data in a packet with $Y=27$.

Unused packets need not be transmitted and incomplete packets shall be filled with the character 0/0.

## PART H: ADDITIONAL EERVICE OPTIONB

## 19. Data for Processing Associated with Displayable Pages

The inclusion of data for processing is signalled by means of the data in Packets with $Y=27$, see Section 12.

### 19.1 Starting Position of Data for Processing

The data for processing on the Page starts with the character code 1/11, which may be in any of the bytes comprising the page. Any bytes before this code are not for processing but may be displayed.

### 19.2 Selection of Protocols

The codes following the $1 / 11$ code determine the decoding Protocol that is in use. There may be any number of codes of the form $3 / x$, followed by any number of codes of the form $2 / x$, followed by a code of the form $7 / x$, where $x$ is in the range 0 to 15.

Note: Character code $1 / 11$ is also used in some national implementations to switch character sets.
20. Television Programme Delivery Control Service

Uses packets with $Y=26$ to carry codes for programming ancillary equipment such as video recorders in order that they respond to the control data as defined in Section 13. These codes can refer to items identifiable by means of a cursor, in the displayable data of the page. Reference should be made to [9].
20.1 Clock Run-In, Framing Code and Packet Address

Bytes 1 to 5 as Sections 9 and 10.

### 20.2 Designation Codes

Byte 6, 4 bits data plus 4 bits Hamming protection.
Codes 0000 to 1111 are sequence labels for up to 16 packets with Y=26 associated with a given page. The sequence is part of that specified in Section 14.6.2. The data groups defined in this Section follow those of Section 14 (see Figure 7). The terminator and CRC check digit complete the sequence as in Sections 14.6.11 and 14.6.12.

The data specified in this Section shall be transmitted as a continuous sequence of packets with $Y=26$ and not interleaved with packets with $Y=26$ providing other functions or with other packets, though other functions may share the first and last packets of such a sequence (see Figure 7).

### 20.3 Data Groups

Bytes 7 to 45, as 13 groups of 3 bytes each. 18 bits data and 6 bits Hamming protection.

```
20.4 Programme Delivery Service
    Facilitates response to data defined in Section 13.3.3 to 13.3.7
    (see also Figures 29 and 30).
20.4.1 Data Group Bit Allocation (See Figure 7)
    6 bits data word A, transmitted least significant bit
    firgt
    5 bits mode description
    7 bits data word B, transmitted least significant bit
    first
    6 bits Hamming protection
20.4.2 Source Definition Data Group
    Data Word A - Country of Origin, see Section 13.3.5.4. 4
    least significant bits specify country of origin. 2 most
    significant bits are set to 1. When set to '0' the group
    is not a Source Definition group and subsequent other data
    groups up to the next Source Definition group shall be
    ignored.
    Mode Description bits set to 01000 invoke the Source
    Definition function (see Figure 9).
    Data Word B - Programme Source: data bits 1 to 6
    correspond respectively to data bits 36 to 31 of Section
    13.3.5.5. Data bit 7 set to '0' corresponds to the first
    set of 64 programme sources, set to '1' corresponds to the
    second set of }64\mathrm{ programme sources.
    In a given page, the Source and Date Definition data
    groups apply to subsequent data groups defining time.
    Either or both of these data groups shall be retransmitted
    when the source and/or date within a given page is to be
    redefined.
```


### 20.4.3 Date Definition Data Group

Data Word A - Month, 2 most significant bits are set to '1'. 4 least significant bits correspond to the equivalent bits in Section 13.3.5.3. Decimal values 49 to 60 define months respectively as January to December. The complete Date Definition Data group and subsequent other data groups up to the next Source or Date Definition group shall be ignored if decimal values 48 and 61 to 63 are transmitted or either of the 2 most significant bits are set to ' 0 '. When set to ' 0 ' the group is not a Date Definition group and subsequent other data groups up to the next Source or Date Definition group shall be ignored.

Mode Description bits set to 01001 invoke the Date Definition function, see Figure 9.

Data Word B - Data bits 1 to 4 specify the day 'units' in BCD form, bit 1 is the least significant bit. Data bits 5 and 6 specify the day 'tens' in $B C D$ form, bit 5 is the least significant bit. Data bit 7 set to ' 0 ', see Figure 7.

The complete Date Definition group and subsequent other data groups up to the next Source or Date Definition group shall be ignored if the following bit values are transmitted:
(a) data bits 1 to 6 corresponding to decimal value 00
(b) data bits 1-4 corresponding to decimal values 13 to 15
(c) data bit 7 set to 1 .

## Local Time Offset

When within a given page, all hours and minutes defined in Section 20.4.5 are in Co-ordinated Universal Time (UTC) this data group need not be transmitted. When all items within a given page have the same Local Time offset it need be transmitted once, immediately following the Date Definition group. When a given page includes items with different Local Time Offsets this group is transmitted before each such corresponding Hours and Minutes definition group.

### 20.4.4.1 Local Time Offset and Cursor Row Definition Group

Data Word A - specifies the row address of the selected cursor position. Decimal values 41 to 63 specify rows 1 to 23 and decimal value 40 specifies row 24. The cursor has no function when located in the header row 0 .

Mode Description bits set to 01100 invoke the Local Time Offset definition function, see Figure 9.

Data Word B - Data bits 1 to 7 specify the Local Time Offset in binary coded quarter hour units.

Bit 1 is the least significant bit and defines $2^{-2}$ hour offset continuing to bit 6 that defines the $2^{3}$ hour offret. Bit 7 defines the sign of the offset, set to ' 1 ' being negative (West of Greenwich) (see Figure 4).

### 20.4.5 Programme Announced Time Data

There shall be a sequence of pairs of data groups, specifying announced starting time hours and minutes for programmes from a common source that run without a break. The announced finishing time hours and minutes shall be included when there is a break in the continuity of the programmes listed and following the final item listed on the page (see Figure 7).

### 20.4.5.1 Announced Starting Time Hours and Cursor Row Definition Group

Data Word A - specifies the row address of the selected cursor position. Decimal values 41 to 63 specify rows 1 to 23 and decimal value 40 specifies row 24. The cursor has no function when located in the header row 0 .

Mode Description bits set to 01010 invoke the Announced Starting Time Hours definition function, see Figure 9.

Data Word B - Data bits 1 to 4 specify the announced starting time hours 'units' in BCD form, bit 1 is the least significant bit. Data bits 5 and 6 specify the announced starting time hours 'tens' in BCD form, bit 5 is the least significant bit. Data bit 7 (Controlled Access Flag - CAF): set to '0', programme item is for free access, set to ' 1 ', programme is for Controlled Access.

The Announced Hours Definition group, the associated Minutes Definition group for Announced Starting and Finishing Time if present shall be ignored if a value of 3 for the hours units digit is exceeded when the hours tens digit has a value of 2 , the value of the hours units exceeds decimal 9, the value for the hours tens exceeds 2. These restrictions do not apply when the month has the value 13 or 14 , in which case the complete label consists of a arbitrary number; or the value 15, when the label has an application-dependent meaning.

### 20.4.5.2 Announced Starting and Finishing Time Minutes and Cursor Position in Row

Data Word A - specifies the character-space within a row occupied by the cursor. Decimal values 0 to 39 specify character-spaces 1 to 40.

Mode Description bits set to 00110 invoke the Announced Starting and Finishing Time Minutes definition function (see Figure 10).

Data Word B - Data bits 1 to 4 specify the minutes units in BCD form, with bit 1 the least significant bit. Data bits 5 to 7 specify the minutes tens in BCD form, with bit 5 the least significant bit.

The complete Announced Hours Definition group, the associated Minutes Definition groups for Announced Starting and Finishing Time if present shall be ignored if a decimal value of the minutes units exceeding 9 or minutes tens exceeding 5, is transmitted. They shall also be ignored if the starting and finishing cursor row position do not correspond. These restrictions do not apply when the month has the value 13 or 14, in which case the complete label consists of a arbitrary number; or the value 15, when the label has an application-dependent meaning.

### 20.4.5.3 Announced Finishing Time Hours, Programme Duration and Cursor Row

Data Word A - specifies the row address of the selected cursor position. Decimal values 41 to 63 specify rows 1 to 23 and decimal value 40 specifies row 24. The cursor has no function when located in the header row 0 .

Row Address Group Mode Description bits set to 01011 invoke the Announced Finishing Time Hours definition function, see Figure 9.

Data Word B - Data bits 1 to 4 specify the Announced Finishing Time Hours 'units' in BCD form, bit $l$ is the least significant bit. Data bits 5 and 6 specify the Announced Finishing Time Hours 'tens' in BCD form, bit 5 is the least significant bit. Data bit 7: set to ' 0 ' announced finishing time is defined as above; set to '1': Programme Duration hours and minutes is defined.

The complete Announced Starting Time Hours Definition group, the associated Minutes Definition group and the Announced Finishing Time Definition groups if present shall be ignored if a value of 3 for the hours units digit is exceeded when the hours tens digit has a value of 2 , the value of the hours units exceeds decimal 9, the value for the hours tens exceeds 2. They ghall also be ignored if the cursor row position in the respective starting and finishing time groups do not correspond. This restriction does not apply when the month has the value 13,14 or 15 , see Section 20.4 .3
20.4.6 Series Code Definition Group

Follows the Programme Timing Data.
Data Word A - set to 110000 defines that the corresponding programme is one of a series. Other Row Address Group values are reserved.

Mode Description bits set to 01101 invoke the Series Code definition function.

Data bits 1 to 7 form a unique code for all programmes in the series, and correspond with the 7 least significant bits of the corresponding data group in the packet 8/30 format 2 in which the most significant bit has the value 'l'.
21. Page-Associated Ancillary Label Data

Uses packets with $Y=25$ to carry a number of LABELS relating to the data in the associated page.

### 21.1 Clock Run-in, Framing Code and Packet Address

Bytes 1 to 5 as sections 9 and 10.

### 21.2 Label Separators

Any Control Character from position 0/0 to 0/7 from Figure 20. Following the standard protocol for packets with $Y=1$ to $Y=24$, Control Character 0/7 (alpha numeric white) is the default starting condition for packets with $Y=25$.

### 21.3 Label Sequence Terminator

Any Control Character from position $1 / 0$ to $1 / 7$ from Figure 20. The interpretation of data in a given packet with $Y=25$, following a terminating control character is reserved.

### 21.4 Label Character Coding

The alpha numeric characters used in the label are those of the default or designated GO character set including the default or specified option. The display colour defined by the separating control character applies to the display of the next label carried by the packet. Control Characters from Figure 20, 0/8 to 0/15 and $1 / 8$ to $1 / 15$ may form part of a label.

## 22. Conditional Access Teletext Service - Page format data

There may be an introductory page of non-scrambled text. When no introductory text for display is required, this page shall at least include the header packet with $Y=0$ and a packet or packets with $Y=27$ to provide links to the conditional access service.

Links are provided in a packet with $Y=27$, designation codes 0100 to 0111 (see Section 12.4). These point to pages with scrambled text and user equipment addressing pages, as identified by data in packets with $\mathrm{Y}=28$ of those pages (see also Figure 31).

### 22.1 Pages with Scrambled Text or Scrambled Data

For the purpose of scrambling, two types of pages are defined:
(a) Pages other than those containing Reformatted Data or Terminal Equipment Addressing Data as in Figure 11. The scrambling process is initialised at the start of each packet. Unused packets need not be transmitted.
(b) Pages containing Reformatted Data, see Figure 11. The scrambling process is initialised at the start of each page.
22.1.1 Pages Not Including Reformatted Data and Not for Terminal Equipment Addressing

The data for transmission in bytes 6 to 45 of packets with $Y=1$ to $Y=25$, plus the 18 data bits in each three byte data group for transmission in packets with $Y=26$, is scrambled, using an enciphering algorithm. The numbers of the packets included in a scrambled text page are specified as in Section 22.2.2.3.

### 22.1.1.1 Parity Protected 7-bit Data

To provide a sequence of complete bytes for scrambling when 7 bit data is used, a most significant bit is added. The resulting bytes are scrambled and the respective bit masked before the odd parity bit added.

### 22.1.1.2 Data in Packets with $Y=26$

To provide a sequence of complete bytes for scrambling, the 18 data bits have 6 most significant bits added. The resulting 3 bytes are scrambled and the respective 6 bits masked before Hamming protection bits are calculated and added.

### 22.1.2 Pages Including Reformatted Data

The data for transmission in bytes 14 to 37 of packets with $Y=0$ and bytes 6 to 45 of packets with $Y=1$ to $Y=25$ are scrambled using an encryption algorithm. The number of bytes included in the page of scrambled data is indicated according to Section 22.2.2.3.
22.1.2.1 Parity Protected 7-bit Data

To provide a sequence of complete bytes for scrambling when 7 bit data is used, a most significant bit is added. The resulting bytes are scrambled and the respective bit masked before the odd parity bit is added.
22.2 Page KeyDescrambling of a scrambled page is by means of a Page Keycontained in a packet with $Y=28$ of a scrambled page.
22.2.1 Page Key Packet - Designation CodeByte 6, 4 bits data plus 4 bits Hamming protection. Databits set to 0010 .
22.2.2 Page Key Packet - Data GroupsBytes 7 to 45 used as 13 groups of 18 bits data plus 6bits Hamming protection See Figure 32(a).
22.2.2.1 Page Key Packet - Page Type Definition Code
This data is not encrypted. The first group of 18 data bits shall be set to designate the page type (see Figure 11).
22.2.2.2 Page Key Packet - 8 Data Groups - Bit Allocation This data is encrypted with Current System Key.



```
22.2.2.3 Page Key Packet - 4 Data Groups
This data is not encrypted.
(1) Continuity Indicator
7 bits Incremented by 1 for each subsequent page with the same service number
(2) Repeat Indicator 1 bit Set to 1 indicates page to be repeated
(3a) For Text Pages
Packet Flags 26 bits
Set to 1 if the packet is present. A flag may indicate a sequence of Packets with \(Y=26\). Least significant packet number is transmitted first.
Set to 0 1 bit
```

오
(3b) For Reformatted Data
Data Length 10 bits
Indicates data length in bytes to be descrambled as a decimal number ( $0=1024$ bytes)
Set to $0 \quad 17$ bits
(4) Scrambling Method

5 bits (see Figure $33(\mathrm{~b})$ )
(5) In Use System Key Label 8 bits
Set to $0 \quad 8$ bits
(6) Cyclic Redundancy Check Word 16 bits (see Figure 8)

For text pages see section 22.1 .1 and $3(a)$
above. The character spaces in packets with $Y=0$ are assumed to contain the character SPACE
(2/0). For Reformatted Data, see Sections 22.1.2 and (3b) above, the check word is calculated over the specified data length.

### 22.3 User Addressing Pages

Access to the Page Key contained in a packet with $Y=28$ of a scrambled page, is by means of the Current and New System Keys. This page contains:
(1) Shared User Data Packets encrypted with the Shared Distribution Key
(2) Unique User Data Packets encrypted with a key that is unique to the user's equipment

System Key Packet encrypted with the New System Rey

The Shared User Data Packets are transmitted relatively frequently and the Unique User Data Packets relatively infrequently.

22.3.1.3 System Rey Packet

This data is encrypted with the New System Key
(1) New System Key 56 bits
(2) Current System Key 56 bits
(3) Set to $0 \quad 16$ bits
22.3.1.4 System Key Packet

This data is not encrypted

| (1) New System Key Label | 8 bits |
| :--- | ---: | :--- |
| (2) Current System Key Label | 8 bits |
| (3) Set to 0 | 56 bits |

### 22.3.2 Shared User Packets - Designation Code

Packet with $Y=1$ to $Y=25$, byte 6.4 bits data plus 4 bits Hamming protection. Data bits set to 0000
22.3.2.1 Shared User Packets - Data Groups

Bytes 7 to 45 used as 13 groups of 18 bits data plus 6 bits Hamming protection. See Figure 34.
22.3.2.2 Shared User Packets - Data Groups - Bit Allocation

This data is not encrypted.
(1) Shared User Address 20 bita
(2) Set to $0 \quad 6$ bits

Followed by data defined in Section 22.3.2.3

```
22.3.2.3 Shared User Packets - Data Groups - Bit
    Allocation
    This data is encrypted with the shared
    Distribution Key, see Section 22.3.3.3
```

| (1) New System Key | 56 bits |
| :--- | ---: |
| (2) User Enabling Bits | 152 bits |
|  |  |
|  |  |

22.3.3 Unique User Packets

Packets with $Y=1$ to $Y=25$. Byte 6, 4 bits data plus 4 bits Hamming protection. Data bits set to 0001.
22.3.3.1 Unique User Packets - Data Groups

Bytes 7 to 45 used as 13 groups of 18 data bits plus 6 bits Hamming protection. See Figure 34.
22.3.3.2 Unique User Packets - Data Groups - Bit Allocation

This data is not encrypted
(1) Unique User Address 32 bits
(2) Set to $0 \quad 2$ bits

Followed by data defined in Section 22.3.3.3

(1) Service Mode (see Section 22.2.2.2) 2 bits
(2) Service Reference Number 2 bits
(3) Current or New Shared Address 20 bits
(4) Unique Equipment Key 56 bits
(5a) Service Modes 1 and 2:
7 Service Numbers of 8 bits 56 bits
(5b) Service Mode 3:
Total credit tokens Purchased for
2 services:
lst Service number 8 bits
Credit Tokens Purchased 20 bits
2nd Service Number 8 bits Credit Tokens Purchased 20 bits
(6) Current or New Shared 56 bits Distribution Key
(7) Current or New User

Enabling Bit Position 8 bits
22.3.4 Initial Page Numbers of Services

Packets with $Y=1$ to $Y=25$. Byte 6, 4 bits data plus 4 bits Hamming protection. Data bits set to 0010 .
22.3.4.1 Service and Associated Initial Page Numbers - Data Groups
Bytes 7 to 45 used as 13 groups of 18 data bitsplus 6 bits Hamming protection. See Figure 34.
22.3.4.2 Service and Associated Initial Page Numbers - Bit Allocation
This data is not encrypted.Each group of 18 data bits includes
(1) Service Number 8 bits
(2) Initial Page Number 8 bits
(3) Set to 0 ..... 2 bits
Any unused groups in a packet are set to 0 .
22.3.5 Independent Data Services - Data Channels and Addresses
Packets with $Y=1$ to $Y=25$. Byte 6, 4 bits data plus 4 bitsHamming protection. Data bits set to 0011.
22.3.5.1 Data Groups
Bytes 7 to 45 used as 13 groups of 18 data bitsplus 6 bits Hamming protection
22.3.5.2 Data Channel Number - Bit Allocation
This data is not encrypted. First group of 18 data bits:

    (1) Data Channel 4 bits see Section 24.3.
    
    (2) Set to 0 14 bits
    22.3.5.3 Service Number and Address - Bit Allocation
This data is not encrypted.
Each pair of Groups of 18 data bits includes:
(1) Service Number 8 bits
(2) Address Length (AL) 3 bits
(3) Address 24 bits
(4) Set to 0 1 bit
Group repeated for each Service Number. Any
unused groups in a packet are set to 0 .

## NOTES TO SECTION 22

## SECURITY OF CONDITIONAL ACCESS TELETEXT SERVICES

The security of the user addressing process may be optimised by the use of "error extension" techniques. Should any bit of the ciphertext have been changed, this causes the message to be totally corrupted when decrypted with the correct key.

In order to provide the property of "error extension" a cipher feedback technique containing a "one-way function" is used with a multiple of encipherments called "rounds". Each round reverses the order of the previous ciphertext bytes, as shown in Figure 35. A typical one-way function having good security is shown in Figure 36.

## Cipher Feedback Algorithm

The secret key $K$ is loaded into the 64 bit key register $K$. The 64 bit key is derived from a 56 bit key that forms the 56 least significant bits. For reformatted data pages and data not in page format, the 8 least significant bits of the 56 are used to look up a corresponding 8 bit value from look-up-tables, LUT in Figure 36. This value is used as the 8 most significant bits of the 64 bit key. For textual page data, the 8 most significant bits are the magazine and row address of each packet forming the page.

The register $R$ is first loaded with a 64 bit secret initial condition $I$, that is constant for the particular security device in the equipment. It is a random number having an impulse autocorrelation function. This data word $I$ is loaded into the register $R$ at the beginning of each round of the encryption or decryption process.

The encryption and decryption processes are represented in Figure 35, the switch being placed in the appropriate positions. The message to be encrypted or decrypted is placed in register $A$ and after the appropriate number of rounds appears in register $B$. The data in register $A$ is taken, byte by byte and the EXCLUSIVE-OR function with the keystream is performed. It is then placed in the $B$ register and the output from the switch is placed in the $R$ register. The previous contents of the $B$ and $R$ registers are shifted along, byte by byte until all the bytes appear in the $B$ register. This process constitutes one round. The next round starts by placing the contents of the $B$ register in the $A$ register but in reverse byte order. At least three rounds are required to produce good ciphertext.

## One Way Function

A suitable one way function is shown in Figure 36. It has a $256 \times 8$ bit look-up-table (LUT) in nine positions. This table contains truly random 'ones' and 'zeros'.

The 64 bit key and the 64 bit contents of register $R$ are added modulo 256 . The resulting 64 bit value is applied to 8 identical look-up tables. A different 1 bit output is taken from each table and these form an 8 bit value. This 8 bit value is applied to a modulo 256 accumulator. This causes each output byte to be influenced by the previous bytes generated during each round of the main algorithm.

The accumulator memory is cleared to zero at the start of each round. The output of the accumulator is applied to a ninth lookup table, identical to the others. Its output forms the key stream of Figure 35 .

## Text Scrambling

The algorithm can also be used to scramble or descramble the user data by placing the switch in Figure 35 in the appropriate position. Only one round is required. To perform either function, the input data is placed in register $A$ and the result appears in register $B$.

## PART I: IMDEPETDEATY DATA 8ERVICES

## 23. Data Transport Protocol

Packets carry information unrelated to, and completely independent of, any accompanying service organised as magazines of pages. See Figure 41.
23.1 Clock Run-In and Framing code

Bytes 1 to 3 as Section 9.
Bytes 4: see Section 23.3.

### 23.2 Designation Code

Byte 5, 4 bits data plus 4 bits Hamming protection. Data bits set to 1111 designate Independent Data Service Packet.

### 23.2.1 Transmission Multiplexing

Television signal data lines carrying these Packets may be included amongst the data lines of a teletext service or may be transmitted using otherwise unused lines. These data lines may always be added at any point in the transmission chain, provided that a new data channel is used (see also Section 23.3).

### 23.3 Data Channel Addressing

Byte 4, 4 bits data plus 4 bits Hamming protection provides 16 Data Channels. These are numbered 0 to 15 and correspond to bit values 0000 to 1111. The least significant bit is transmitted first.

For the services defined in Section 23 four data channels are allocated:
Data Channel Number Address Bit Values in Transmission Order 0001
9 1001
10
0101
11
1101
Note that data channel 0 , defined by the data bits being set to 0000, is used for the Broadcasting Service Data Packet, see Section 13.
23.4 Format Type (FT)

Byte 6, 4 bits data plus 4 bits Hamming protection. Format Type $A$, defined by:
Bit 1...get to 0
Bit 2...set to 1: repeat packet facility applies, see
Section 23.4.3
Bit 2...set to 0 : no repeat facility
Bit 3...set to 1: explicit continuity indicator included
Bit 3...set to 0 : continuity indicator is implicit
Bit 4...set to 1: Data Length Byte in Use
Bit 4...set to 0: Data Length Byte not in use
23.4.1 Service Packet - Interpretation and Address Length (IAL)

Byte 7, 4 bits data plus 4 bits Hamming protection. The first three bits define the number of immediately following Hamming coded bytes which are allocated to defining the service packet address. All three bits set to 0 indicate that there is no service packet address within the data line. Each increment in binary value adds 4 bits to the address length up to a maximum of 24 bits. All three bits set to 1 is reserved for future extensions.

The fourth bit set to 0 defines data as independent of the contents of any other channel or address. The fourth bit set to 1 indicates that interpretation of the data may require the use of data in other channels or with other addresses as defined by the application.

### 23.4.2 Service Packet Addresses

When present bytes 8 to 13, see Section 23.4.1.
NOTE: When differentiated by the appropriate address length data in byte 7 , the less significant bytes of an address may constitute another complete address in the same data channel. Thus, for example, the 24 -bit address ABC123 can co-exist with the 20 -bit address BC123 and the 8 -bit address 23.
23.4.3 Service Packet - Repeat Indicator (RI)

This byte follows the Service Packet Address and is only present when the Format Type bits are appropriately set, see Section 23.4. The first 4 bits are set to 0 when a new packet of that service data channel is first transmitted and shall be incremented modulo-16 on subsequent repeats.

The next three bits are reserved for future extensions. The last bit is set to 0 to indicate that no further repeats of the current packet should be expected. This last bit shall be set to 1 when a further repeat is to be expected.

### 23.4.4 Packet Continuity Indicator (CI)

This byte follows the service packet address or the Service Packet Repeat Indicator, if present. It is only present when the Format Type bits are appropriately set, see section 23.4. It represents an 8-bit number which is incremented modulo-256 with each new packet of the same address on the same data channel. It is not incremented on repeated transmissions of the same packet.

### 23.4.5 Data Length Byte (DL)

This byte follows the Service Packet Address or the Service Packet Repeat Indicator or Packet Continuity Indicator if they are present. It is only present when the Format Type bits are appropriately set, see Section 23.4.

The two most significant bits are not defined. The remaining six bits define the number of 8-bit bytes of user data intended to be delivered to the user. The count is taken from the start of the User Data Byte Group and includes any dummy bytes, see Section 23.4.6.

The DL byte is included when it is necessary to send an incompletely filled packet. Any remaining bytes of the user data group are not defined but are subject to the CRC, see Section 23.4.7.

The six data bits defining the data byte length may be set to 0 , to keep a data service channel open when there is no data for delivery to the user.

## User Data Group

The remaining data bytes in the data line, excepting the last two, constitute the data carried for users of the service bearing that Service Packet Address on that data channel. The number of bytes available depends upon the address length, whether the repeat facility is used, and whether the continuity indicator is implicit or explicit. Thus there are between 28 and 36 data bytes available.

Certain forms of coding may give rise to long strings of $0 s$ or 18. It is desirable to remove these from the transmitted data field to ensure reliable operation of all equipment that may process the signal. When within any user data group a sequence of eight consecutive bytes containing all 0 os or eight consecutive bytes containing all ls occurs, taken together with its CI byte if present, a following dummy byte will be inserted. This dummy byte is included in the calculation of the CRC, (see Section 23.4.7) but is otherwise ignored by the decoder.

NOTE: Decoders must be designed to recognise these dummy bytes. However their inclusion may in the future no longer be necessary and it would be desirable to omit them to increase efficiency. It is therefore recommended that decoders should be capable of convenient modification or adjustment when this occurs.

### 23.4.7 Cyclic Redundancy Check Word

The last two bytes contain a Cyclic Redundancy Check on the user data group (see Section 23.4.6) and on any Continuity indicator (CI) or Data Length (DL) byte if present (see Sections 23.4.4 and 23.4.5).

### 23.4.7.1 Check Word Generation

The data to be checked is considered as a polynomial in $x$ with the highest degree term transmitted first and the term of degree zero last. This is divided, using modulo-2 arithmetic by the polynomial:

$$
x^{16}+x^{9}+x^{7}+x^{4}+x^{0}
$$

The remainder from this process, with the highest term transmitted first, is the CRC.

When an implicit continuity indicator is signalled by the third message bit of the FT byte, the transmitted CRC is modified such that the described generation process results in the register containing the 8-bit continuity indicator byte twice, with the least significant bit at the right-hand end.

### 23.4.7.2 Check Result

The register of Figure 42 is set to $0 s$. The serial data followed by the CRC is then entered. The check is satisfied if the register again contains all $0 s$.

## NOTE CONCERNING TRANSMISSION SEQUENCE:

For any Service Packet Address, the corresponding serial data stream is divided into User Data Groups. These must be transmitted in the correct sequence, which is monitored by the continuity indicator. Provided that the Repeat Indicator is used, each group may be repeated any number of times before the next is sent. There may or may not be an interval between consecutive data-lines with the same Service Packet Address.

Data lines carrying different Service Packet Addresses may be combined in any order to form a Data channel, provided that the sequence for each Service Packet Address is not disturbed.

Data-lines from different Data Channels may be combined in any order provided that the sequence within each contributing source is not disturbed.

## 24. Conditional Access for Block Mode Data

Data in the User Data Group (see Section 23.4.6) of a number of Independent Data Service Packets are linked to form Data Blocks. They may contain Messages concerning access that are not Encrypted, Encrypted Messages concerning access and User Data to be communicated.

### 24.1 Block Separator

Byte 1 of a data block.
Blocks of data are separated by the transmission of the character hex 10 (DLE). If the character DLE occurs within the Data Block it shall be repeated to indicate that it is not a block separator.

### 24.2 Block Formats

Two alternative formats are defined. Format $A$ is defined in Section 24.3 and Format $B$ in Section 24.4.

### 24.3 Block Format A

"User Data" is the data stream originated by a sender and intended for delivery to a specified recipient or group of recipients. "Messages" are groups of data concerned with the access to User Data and are for controlling the decoder or placing it in operation. A suitable encryption algorithm is shown in Figure 35.

### 24.3.1 Block Type

Byte 2 of a data block defines the block type. This byte is of the form $04 / 0 x$ where $x$ may have a value 0 : $0,1,2,3,4,5,6$ or 7 .

Block Type Code Interpretation
04/00 Block contains User Data not
Scrambled; see Figure 37(a)
04/01

04/02
Block contains Key Message not Encrypted, and the Encrypted Key Message
Service Numbers and Sequence
Numbers, not encrypted and
Scrambled User Data.
04/03 System Key Message Block
04/04
Shared User Message Block
04/05
04/06
$04 / 07$
Unique User Message Block
Service Address Message Block for
Independent Data Services
Service Address Message Block for
Page Format Services
24.3.2 Primary Block Key Messages

Block Type Code 04/01. See Figure 37(b).
24.3.2.1 Block Rey Message Bit Allocation

This data is not encrypted
In Use System Key Label 8 bits
followed by data defined in Section 24.3.2.2.
24.3.2.2 Block Key Message - Bit Allocation

This data is encrypted with the Current System Key

| (1) Service Modes |  |  |
| :--- | :--- | :--- |
| Bit 2 | Bit 1 | 2 bits |
| 0 | 0 | 256 Services Non-Tiered |
| 0 | 1 | 64 Services Tiered <br> 1 |
|  | 0 | 256 Services with Credit <br> Tokens |
| 1 | 1 | Not Assigned |


| (2) Set to 0 | 6 bits |
| :--- | :--- | :--- |
| (3) Current System Rey | 56 bits |
| (4) Service Identification Number | 8 bits |



### 24.3.4.2 System-Key Message Block Bit Allocation

This data is encrypted with the New System Key
(1) New System Key 56 bits
(2) Current System Key 56 bits
24.3.5 Shared-User Message Block

Block Type Code 04/04. See Figure 37(e).
24.3.5.1 Shared-User Message Block - Bit Allocation

This data is not encrypted
(1) Shared User Address 20 bits
(2) Set to 0 data defined in section 24.3 .5 bits

Followed by data defined in Section 24.3.5.2.
24.3.5.2 Shared-User Message Block - Bit Allocation

This data is encrypted with the Shared
Distribution key, see Section 24.3.6.2.
(1) New System Key

56 bits
(2) User Enabling Bits 152 bits $0=$ disabled; $1=e n a b l e d$
24.3.6 Unique-User Message Block

Block Type Code 04/05. See Figure 37(f).
24.3.6.1 Unique-User Message Block - Bit Allocation

This data is not encrypted Unique User Address 32 bits Followed by data defined in Section 24.3.6.2.


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(5b) Service Mode 3: Total credit tokens purchased for 2 services: lat Service Number 8 bits Credit Tokens Purchased 20 bits 2nd Service Number 8 bits Credit Tokens Purchased 20 bits
(6) Current or New Shared 56 bits Distribution Key (7) Current or New User 8 bits
24.3.7 Service Address Message Block - Independent Data Service Block Type Code 04/06. See Figure 37(g)
24.3.7.1 Service Address Message Block - Bit Allocation
This data is not encrypted
(1) Service Number: 8 bits
(2) Address Length (AL) as Section 3 bits
23.4.1.
(3) Set to $0 \quad 5$ bits
(4) Service Address up to 24 bits
This group is repeated for each service number
24.3.8 Service Address Message Block - Page Format Service
Block Type Code 04/07. See Figure 37(h)
24.3.8.1 Service Address Message Block - Magazine Number
- Bit Allocation
This data is not encrypted
Byte following Block Type Code
(1) Magazine Number 3 bits
(2) Set to $0 \quad 5$ bits
This byte is not repeated for each service
number.
24.3.8.2 Service Address Message Block - Initial Page
Number - Bit Allocation
This data is not encrypted
(1) Service Number 8 bits
(2) Initial Page Number 8 bits
This group is repeated for each service number.

### 24.4 Block Format B

User Data is the data stream originated by a sender and intended for delivery to a specified recipient or group of recipients, see Figure 38. Messages are groups of data concerned with the access to User Data and are for controlling the decoder or placing it in operation, see Figure 38. A suitable encryption algorithm is shown in Figure 39, see also the Notes to Section 24.

### 24.4.1 Block Type

Byte 2 of a data block defines the block type. This byte is of the form 5/X where $X$ may have the hexadecimal value of $0,7,8,9, A, B, C, D, E$ or $F$.

Block Type Code Interpretation
5/0 Non-scrambled Channel and Terminator
5/7 Common User-Data
5/8 User-Data Key Message
5/9 Key-Conversion Message
5/A Shared System-Key Message
5/B Group User-Data
5/C Unique System-Key Message
5/D Shared Equipment-Key Message
5/E Over-Air-Credit Message
5/F Unique User-Data
Groups of bytes in Messages are transmitted with the least significant byte first. Groups of bits in Messages defining a Mask are transmitted with mask bit 0 first.
24.4.2 Non-Scrambled Channel

Bytes following Block Type Code 5/0 up to the next block separator (DLE) do not include scrambled or encrypted data. This block type is used to terminate a previous block, when there is no other data for inclusion.
24.4.3 Common User Data Block

Block Type Code $5 / 7$, see Figure $40(\mathrm{~b})$.

```
24.4.3.1 Common User Data Block - Bit Allocation
This data is not encrypted
    In Use System Key Label 1 byte
    Cipher Initial Variable 1 byte
```


### 24.4.3.2 Common User Data Block - User Data <br> Any number of bytes scrambled by cipher stream using method 2 in Output Feedback Mode, see Figure $33(\mathrm{~b})$.

```
24.4.4 Group User Data Block
    Block Type Code 5/B, see Figure 40(c).
    24.4.4.1 Group User-Data Block - Bit Allocation
        This data is not encrypted
        Shared User Address 3 bytes
        Cipher Initial Variable
            1 byte
    24.4.4.2 Group User-Data Block - User Data
    Any number of bytes scrambled by cipher stream
        using method 2 in Output Feedback Mode, see
        Figure 33(b).
24.4.5 Unique User-Data Block
    Block Type Code 5/F, see Figure 40(d).
    24.4.5.1 Unique User-Data Block - Bit Allocation
    This data is not encrypted
    Unique User Address 5 bytes
    Cipher Initial Variable 1 byte
    24.4.5.2 Unique User-Data Block - User Data - Bit
    Allocation
    Any number of bytes scrambled by cipher stream
    using method 2 in Output Feedback Mode, see
    Figure 33(b).
```


### 24.4.6 User-Data Key Message Block

```
Block Type Code \(5 / 8\), see Figure \(40(\mathrm{e})\).
24.4.6.1 User-Data Key Message Block - Bit Allocation
This data is not encrypted Decoder Secure Module Address 1 byte In Use System Key Label 1 byte User Data Key Label 1 byte
```

```
24.4.6.2 User-Data Key Message Block - Bit Allocation
```

24.4.6.2 User-Data Key Message Block - Bit Allocation
This data is encrypted with current system key
This data is encrypted with current system key
using encryption method 2 in Differential Code
using encryption method 2 in Differential Code
Book Mode, see Figure 33(a).
Book Mode, see Figure 33(a).
Part Current System Key 5 least significant
Part Current System Key 5 least significant
bytes
bytes
Sub-service 1 byte
Sub-service 1 byte
Item Value:
Item Value:
fractional part of value: l byte
fractional part of value: l byte
whole unit part of value: }1\mathrm{ byte
whole unit part of value: }1\mathrm{ byte
User Data Key: }8\mathrm{ bytes

```
    User Data Key: }8\mathrm{ bytes
```

```
24.4.7 Key-Conversion Message Block
    Block Type Code 5/9, see Figure 40(f).
    24.4.7.1 Key-Conversion Message Block - Bit Allocation
    This data is not encrypted
    New System Key Label l byte
    Current System Rey Label l byte
    24.4.7.2 Key-Conversion Message Block - Bit Allocation 
24.4.8 Shared System-Key - Message Block
    Block Type Code 5/A, see Figure 40(g).
    24.4.8.1 Shared System-Key - Message Block - Bit
    Allocation
    This data is not encrypted
    Shared User Address 3 bytes
24.4.8.2 Shared System-Key - Message Block - Bit
    Allocation
    This data is encrypted with the Shared Equipment
    Key using encryption method 2 in Differential
    Code Book Mode, see Figure 33(a).
    Part Shared Equipment Rey: }7\mathrm{ least
                                    significant bytes
    In Use System Key Label
                            l byte
    User Enabling Mask: 32 bytes
    In Use System Key: }8\mathrm{ bytes
24.4.9 Unique System-Key Message Block
Block Type Code 5/C, see Figure 40(h).
```

24.4.9.1 Unique System-Key - Message Block - Bit Allocation

```
    This data is not encrypted
    Unique User Address 5 bytes
24.4.9.2 Unique System-Key Message Block - Bit Allocation
    This data is encrypted with the Unique Equipment
    Key using encryption method 2 in Differential
    Code Book Mode, see Figure 33(a).
    Part Unique Equipment Key: 7 least
                                    significant bytes
    Current System Key Label:
    Sub-service Mask:
    l byte
    l6 bytes
    Current System Key: }8\mathrm{ bytes
```


24.4.11 Over-air Credit Message Block

Block Type Code 5/E, see Figure 40(k).
24.4.11.1 Over-air Credit Message Block - Bit Allocation

This data is not encrypted Unique User Address 5 bytes
24.4.11.2 Over-air Credit Message Block - Bit Allocation

This data is encrypted with the unique equipment key using encryption method 2 in Differential Code Book Mode, see Figure $33(\mathrm{a})$. Previous Total Credit: 4 bytes Unique Equipment Key: 8 bytes New Total Credit: 4 bytes

## NOTES TO SECTION 24

## SIMPLE BLOCR ENCIPHERMENT ALGORITHM

## Encryption Method 2 Figure $33(a)$

Current Standardization activity on data encipherment is based on the use of a 64-bit cipher in one of four modes. The key word length is at least 48 bits.

An encipherment algorithm is proposed which is relatively simple to implement in software or hardware, including VSLI, with adequate security for commercial broadcasting applications. It uses 64-bit blocks and a 48-bit key, extendable to 64-bits.

The operation uses a device, called a permuter, which under the command of a 2-bit control input, converts an 8-bit input word to an 8-bit output word. The conversion is by means of one of four stored permutations of the 256 input states to 256 output states. Because these are permutations, not random 'look-up tables', the operation is reversible, each output state corresponding to one and only one input state for a given control word. Thus for every permuter there is a complementary permuter which acts in the reverse direction.

A convenient embodiment of the permuter is a 1024 by 8-bit ROM. It has a 10 bit input, 2 bits being assigned to the control function and the remainder are the 8-bit input word. 24 permuters are required for full-speed operation. It is recommended that for commercial broadcasting applications, all the permuters are identical but the method describes the general case where they can all be different.

In Figure 39 the permuters are shown as 8-by-8 squares with 8 inputs on the left-hand edge and 8 outputs on the right-hand edge. They are in banks of 8 so that they handle 64 bits at once. The ordered 64 -bit input is applied to the 64 input points using a fixed but irregular (e.g.random permutation of the 64 items) distribution. The output of the first bank of permuters (a-h) forms the input to the second bank (j-r). It is desirable that each permuter in the second bank has one input from each of the permuters in the first bank. The same consideration applies when coupling the second bank to the third bank of permuters (s-z). This is shown in Figure 39 by simple transposition. Although a more random fixed permutation would be of advantage, it is not considered necessary for this application. The output 64 bits are assembled using another fixed but irregular distribution. The 48 -bit key is applied to the 24 permuters such that each receives 2 unique control bits in a fixed and not necessarily irregular way. The principles can be readily extended to a 64 bit key by adding another bank of permuters, using the same pattern.

By reading Figure 39 from right to left, with complementary permuters, but using the original control bits from the key word, the reverse process takes place to decipher the 64-bit block. The same key is used as that for the original encipherment.

A number of equivalent terms are used in current descriptions of controlled access systems and the encipherment of data. In Sections 22 and 24 a practical degree of uniformity has been employed in respect of the range of formats. However the following notes will assist when relating these sections to other documents and standards.

1. In broadcasting applications generally the term "Scrambling" is used for the operation of rendering programme or other material for transmission of no value to an unauthorised recipient. The complementary process is "Descrambling". This meaning is used in this specification. The terms encryption/decryption and encipherment/decipherment are in use for aspects of these functions.
2. In broadcasting applications generally, the term "Encryption" is used for the operation of rendering a "Descrambling Key" (i.e. a data word) only accessible to authorised recipients. The complementary process is "Decryption".
3. The term "Encipherment" and its complement "Decipherment" is used for the process of using a cipher (i.e a Key) to carry out the processes o: Scrambling/Descrambling and Encryption/Decryption.
4. The term "System Key" is used in this specification. An alternative term is "Period Key".
5. It is necessary to define and label transmitted keys as being "Current' or "New". Alternative terms are respectively "Result" and "Source".
6. The Electronic Code Book (ECB) and the Output Feed Back technique are the encipherment modes for a 64-bit block cipher as defined in [11].

## PART J: APPEADICE8

## APPENDII 1

RECOMMENDED ORDER OF TRANSMISSION FOR PACKETS WITH $Y=26, Y=27$ AND $Y=28$
To permit most efficient operation of the decoder it is recommended that following the transmission of a page header packet $(Y=0)$, any packets with $Y=27, Y=28$ and $Y=26$ are transmitted, in that order.

If a further header packet $(Y=0)$ is transmitted, following the transmission of packets with $Y=27, Y=28$ and $Y=26$, then those packets must be re-transmitted.

## APPEINII 2

## DECODER MEMORY BEHAVIOUR

Although this specification does not explicitly introduce any limitations based on the expected size or speed of the decoder memory, certain decoders may have limitations in these respects as follows:

1. The number of attributes applicable to any one display row may be limited to 40.
2. The DRCS memory may be limited to $16 k$ bita.
3. Up to 20 ms may be required to erase the decoder memory after the receipt of a page header packet ( $Y=0$ ). Thus it may be necessary for this period to elapse, after transmission of a page header packet before more data relating to that page is transmitted.

## APPENDIX 3

## HAMMING PROTECTED DATA

Certain data bytes are protected by Hamming Codes, two forms are used:

1. Hamming ( 8,4 ) - Bytes with four data bits and four protection bits.

In this form bits b1, b3, b5 and b7 are the protection bits; bits b2, b4, b6 and b8 carry the data.

The Parity Checks Pn are:

| P1 | $b 8, b 6, b 2, b 1$ |
| :--- | :--- |
| P2 | $b 8, b 4, b 3, b 2$ |
| P3 | $b 6, b 5$ b4, b2 |
| P4 bits 1 to 8 inclusive |  |

2. Hamming (24,18) - Groups of three bytes with 18 data bits and 6 protection bits.

In this form bits b1, b2, b4, b8, b16, b24 are the protection bits; bits b3, b5, b6, b7, b9, b10, b11, b12, b13, b14, b15, b17, b18, b19, b20, b21, b22, b23 carry the data.

The Parity Checks Pn are:
Pl all odd numbered bits 1 to 23 inclusive
P2 b2, b3, b6, b7, b10, b11, b14, b15, b18, b19, b22, b23
P3 b4, b5, b6, b7, b12, b13, b14, b15, b20, b21, b22, b23
P4 bits numbered b8 to bl5 inclusive
P5 bits numbered b16 to b23 inclusive
P6 bits numbered b1 to b24 inclusive

## APPENDIX 4

## CODES OF PRACTICE FOR OPTIONAL USER-FRIENDLY PAGE ACCESS METHODS

METHOD 1 - "FLOF"

## A. INTRODUCTION

The object of this code of Practice is to facilitate the addition of new service features, the design of receivers to exploit them and the introduction of the concepts to the market.

The intention is to provide ease of page selection, comprehensive use of the data base and a reduction in page acquisition time.

This Code of Practice is compatible with the EBU TELETEXT SPECIFICATION but is not a part of that specification. It includes information applicable to broadcasters and receiver designers.

## B. OBJECTIVES

1. Simple User Control
2. Easier Selection of Information
3. Reduced Page Access Time
4. Compatibility with Existing Products in the Field
5. Compatibility with Existing Data Base Structures
6. Minimum Transmission Overheads
7. Moderate Editorial Overheads
8. Optimum Cost Effectiveness
9. Flexibility in Use
10. Unambiguous Operation
11. Compatibility with Decoders Having Storage Facilities for Any Number of Pages

## C. TRANSMITTED PACKETS

In addition to the packets with $Y=0$ to $Y=23$, used for level one display, the service features included in this code of practice require the following extension packets:

Packets with Y=24 Page Data Extension Packet
In this application, this packet contains user 'prompt' information for display, provided by the editor. It associates a given key on the user's control unit with a linked page address included in packets with $Y=27$, gee Section 12.2.

Packets with Y=27 Designation Code 0000 Page Service Data Packet
This packet contains the addresses of linked pages, link control information, display row 24 flag and error detection data, see section 12.2.

Packet 8/30 Television Service Data Packet
Format 1 of this packet applies and contains the address of the 'Initial Teletext Page', 'Status Message', 'Time Offset from Universal Co-ordinated Time','Universal Co-ordinated Time', 'Modified Julian Date' and 'Television Programme Group Label'. When a given category of the above data is not included, fall-back codes will be included.

## D. EXTENSION PACKET APPLICATION

## D. 1 Packet 8/30

| D.1.1 | The data defined in Sections 13.1 to 13.2 .9 shall transmitted. |
| :---: | :---: |
| D. 1.2 | Decoders intended for use according to this code of practice shall process the designation code (Section 13.2.1), recognising multiplexed or non multiplexed functions. |
| D. 1.3 | Decoders intended for use according to this code of practice shall acquire the initial teletext page on switch-on or changing channel, see Section 13.2.2. |
| D. 1.4 | Decoders intended for use according to this code of practice may optionally process these data items: |
|  | - Network Identification (Section 13.2.3) <br> - Time Offset Code (Section 13.2.4) <br> - Modified Julian Date (Section 13.2.5) <br> - Co-ordinated Universal Time (Section 13.2.6) |
| D. 1.5 | Television Programe Label data (Sections 13.2.7 and 13.2.8) Decoders intended for use according to this code of practice are not required to process these data items. |
| D. 1.6 | Status Display (Section 13.2.9) Decoders intended for use with according to this code of practice may display the status message on switch-on and on changing channel. It is not intended to be presented in the teletext mode but is inserted in a box in the normal television picture display. start Box and End Box control characters are not transmitted. The box must therefore be generated by the decoder. The display of the message shall be terminated after a suitable period. User controls may be provided to reactivate the display for a further period and to deactivate the status display facility. |

## D. 2 Packets with Y=27 Automatic Page Acquisition

A decoder intended for use according to this code of practice shall process the linked page addresses (Section 12.2 ) to acquire pages automatically. The maximum number of pages stored is limited by the decoder memory capacity.
D.2.1 Coloured Keys - Prompt ModeA decoder intended for use according to this code ofpractice will provide four keys on the user's control unit,associated with the first four linked page addresses, intransmission order. These keys shall be respectivelycoloured red, green, yellow and cyan and relate to theassociated displayable prompts that may be transmitted in apacket with $Y=24$.
D. 2.2 Link Address Number 4The fifth link address is not used in the functions definedabove but may be used for additional Automatic PageAcquisition.
D. 2.3 Link Address Number 5
A decoder intended for use with this code of practice willprovide an INDEX KEY on the user's control unit. The sixthlink address is associated with this key and the linked pageshall carry suitable introductory material to assist theuser. The use of the Index Key is not associated with aprompt transmitted in a packet with $Y=24$.
D.2.4 Display Row 24 Flag (Section 12.2.3.2)
Data bit 4 set to 1 :
Packet with $Y=24$, if present, is to be displayed at thebottom of the associated page. Such a packet 24 may includethe prompts as defined above.
Data bit 4 set to 0:
Packet with $Y=24$ is not to be displayed but may be presentfor other applications.
(The functions of data bits 1,2 and 3 of the Link Control Data are as defined in Section 12.2.3.1 and are not part of this code of practice).
Decoders compatible with this code of practice shall process the Display Row 24 Flag bit 4 as described above.
D.2.5 Basic Page Check Word (Section 12.3)
Decoders intended for use according to this code of practice may optionally process this check word.
D. 3 Packets with $Y=24$

This packet is used in this code of practice to assist the user in page selection. The relationship between the displayed prompts and the users control keys should be clear.

Up to four prompts can be included in this packet, respectively distinguished by the inclusion of the colours red, green, yellow and cyan. The display should be such that a prompt can be distinguished by position as well as by colour.

The number of characters used in the prompts is only limited by the availability of 40 character codes in packets with $Y=24$.

The relationship between prompts transmitted in a packet with $Y=24$ and the link addresses in packets with $Y=27$ is as defined in Section D.2.1 and D.2.4.

## E. SERVICE COMBINATIONS

The following fall-back conditions are defined:

## E. 1 No Packet 8/30

Multiplexed operation is assumed, with an initial page address 100, page sub-code being interpreted as 'don't care'. No status message is to be displayed.
E. 2 No Packet with $Y=27$ and No Packet with $Y=24$

The index key will cause the initial page to be displayed; the coloured keys will be disabled.
E. 3 Packet with $Y=27$ but No Packet with $Y=24$

At least an index link in the appropriate link number shall be provided.
E. 4 A Packet with $Y=24$ included, other than for prompts in conformity with this code of practice

A packet with $Y=24$ may be included in any sequence of data. It shall only be displayed when the Display Row 24 Flag Bit is set to 1 , see Section D.2.4.

## F. CONFORMANCE

F.1 To conform with this code of practice the transmission shall include: - a packet 8/30 including the data specified in Section 13.2; and a packet with $Y=27$ including the data specified in Section 12.2; and a packet with $Y=24$ including data as specified in D.3.

Note: The inclusion of prompts in a packet with $Y=24$ of a given page is under editorial control since some material may not be suitable, e.g. due to rapidly changing content.
F. 2 To conform with this code of practice a decoder shall:

- in addition to packets with $Y=0$ to $Y=23$, at least acquire packets with $Y=24, Y=27$ and $8 / 30$;
and
at least process the data in packet $8 / 30$ as specified in Section D.1.2 and D.1.3;
and
- at least process the data in packets with $Y=27$ as specified in Sections D. 2 but excepting Sections D.2.2 and D.2.5.
and
at least display the data in the packet $Y=24$ as specified in Section D. 3.


## G. ADDITIONAL DECODER FEATURES

Certain decoder features have been identified that do not require the transmission of additional data and do not need to be taken into account in the editorial processes. They do not form part of this code of practice but are included for the guidance of equipment designers.

## G. 1 Decoder Display When there is No Transmitted Packet with Y=24

When the link addresses in a packet with $Y=27$ are not associated with prompts and no packet with $Y=24$ is transmitted, the decoder may generate a display. It is recommended that this may include the magazine and page numbers of the first four link addresses (numbers 0 to 3) in a packet with $Y=27$. Any page sub-code should not be displayed.

To utilise this facility, the link address number 0 , displayed in red and relating to the red key, may be used editorially defining the 'next logical page'; the link number 1 , displayed in green and relating to the green key, may be used editorially defining a 'browse' among various pages. The position and mode of this display is optional but the relationship to the keys should be clear.
G. 2 Previous Page

It is recommended that a 'Previous Page' key be included in the user controls. This key should toggle between the currently selected page and the page previously displayed.
G. 3 Sub-title Key

It is recommended that a key be included to provide simple access to the sub-title service.

## G. 4 Other Methods of Page Selection

Notwithstanding the inclusion of the data defined in this code of practice in the transmission and the inclusion of the related keys on a user's control unit, other methods of page selection may still be provided. While this code of practice includes a range of features appropriate to an enhanced service, it does not preclude other functions, either included in the transmission or in the decoder, that provide further features.

METHOD 2 - "TOP"

## A. INTRODUCTION

The "TOP" (Table of Pages) System specified in [11] and described below is an addition to the teletext service which can be offered to the teletext user by the service provider. A user-friendly page access such a teletext transmission cycle can be arranged of suitably-equipped decoders. The selection of the pages is accomplished by special function keys on the user's teletext television remote-control unit. TOP also facilitates users rapid orientation in the teletext service structure together with an improved user guidance.

The TOP System Code of Practice is compatible with the EBU TELETEXT SPECIFICATION but is not part of that specification. It includes information applicable to broadcasters and receiver designers.

## B. OBJECTIVES

The TOP system is intended to fulfill the following objectives

1. Simple user control
2. Rapid orientation with in the teletext service structure
3. Reduced page access time
4. Better use of the decoder data base
5. Direct page access facilities
6. Minimum transmission overhead
7. Moderate editorial overhead
8. Low additional costs for providers and viewers
9. Compatible with multipage decoders
10. Additional features improving user's guidance

## C. MAIN PRINCIPLE

The basic idea underlying the TOP System is to arrange the pages within a teletext transmission cycle according to specific themes comparable to the structure of themes established by ordinary file-card boxes. All "news" pages, for example, are assigned to a specific file-card box; further "blocks" of pages are conceivable: one for "sport", one for "service", one for "TV programme preview" and so on. Correspondingly, as it is also customary for file-card boxes, a further sub-division is provided: each "block" is composed of different "groups", for instance, the "news" block might be subdivided into the groups "politics", "economics" and "culture".

- User-controlled selection of pages

In accordance with the partition into blocks and groups, the selection of pages is established by way of four specially-marked keys on the decoder remote-control unit:

| lst key $(n+"): \quad$ leads to the next page in the "file-card" box |  |
| :--- | :--- |
| 2nd key (e.g. "red"): leads to the first page of the next group |  |
| 3rd key (e.g. "blue"): leads to the next block and |  |
| 4th key ("-"): | leads back to the page(s) last seen. |

In the commentary row (25th text row) at the bottom of the display screen, two fields, in the same colour as used by the corresponding keys, serve as guide for the user. This commentary row generated by the decoder announces the title of the block just selected, the title of the next group in that block and the title of the next block. In addition, the IV programme-preview pages can directly be selected by pressing a special key on the remote-control unit.

## - Direct page access

In place of a step-by-step keying-in from one block to the next and from one group to the next, direct page access may also be offered by the TOP System. By way of a "guide" page listing all the blocks together with their groups a user can mark the desired group within that page to have direct access to the pages of this group.

In addition, specific pages can be marked by the teletext editor in order to generate a special "review" page for direct access by the viewer. Thus the editor can provide a simple and comfortable access for pages, which are particularly important or frequently viewed by the viewers.

## - Additional features

By evaluating the TOP tables included in the teletext transmission cycle additional announcements such as "Page not included", "Please wait", "Multipage with 4 sub-pages" can be issued by the decoder in the commentary row without having to await the appearance of the page demanded.

By use of TOP the waiting time for requested pages may also be considerably reduced. Where an appropriate design of multi-page memory decoder is used. Stored pages in the vicinity of a selected page which have a high probability of viewer interest within a group or a block may be stored while the viewer is reading the last-selected page.

TOP contributes, to providing user-friendly guidance through the multitude of pages within a complex-structured teletext programme, and also to optimally exploitating of the storage capacity offered by the teletext decoder.

## D. IMPLEMENTATION

TOP is based on the inclusion of four kinds of "Table" pages within the teletext transmission cycle:

- the BASIC TOP TABLE,
- the PAGE-LINKING TABLE,
- the MULTI-PAGE TABLE and
- the ADDITIONAL INFORMATION TABLE.


## D. 1 Basic TOP table

In the BASIC TOP TABLES all teletext pages included in the transmission cycle are listed together with special flags set by the teletext editor for identifying blocks, groups, programmepreview pages and subtitles.
D.1.1 Structure

All pages with decimal page numbers admissible in the teletext system (from 100 to 899) are placed in the BASIC TOP TABLE, each with a distinct position whitin a coordinate system; thus, for example, the position for page 100 is allocated to column 0 of row 1 .

The relation between the page number and the position coordinates in the Basic TOP Table is according to the following formula:

$$
N=100+C+40 \times(R-1)
$$

where $\quad N=$ decimal page number (100...899)
$R=$ row address $C=$ column address
(0...39)
D.1. 2 Coding

The transmission codes for the BASIC TOP TABLE are formed from the source codes by applying the $(8,4)$ Hamming code.

The composition of the page sub-code of the BASIC TOP TABLE is as follows:


The seven most-significant bits of the subcode are permanently set to logical "1"; accordingly, the page sub-code can have values from $3 F 00$ to $3 F 7 F$.

## D. 2 Page-linking table

The PAGE-LINKING TABLE contains a list of page numbers on which other tables of the TOP system are transmitted. In addition to the page numbers, the PAGE-LINKING TABLE also contains a list of the various TOP tables types MULTI-PAGE TABLE or ADDITIONAL INFORMATION TABLE).

## D.2.1 Structure

The PAGE-LINKING TABLE is part of the teletext page carrying the BASIC TOP TABLE. Use is made of rows 21 and 22 which are not utilized by the BASIC TOP TABLE. Up to five page numbers and the corresponding table indications may be given in each of these two rows. Altogether, it is possible to address $2 \times 5=10$ additional TOP tables in row 21 and 22.

## D.2.2 Coding

The coding of the page numbers together with the table identification is according to [11]. The transmission codes are formed from the source codes by applying the $(8,4)$ Hamming code. If fewer than ten table-page numbers are entered into the table, the replacement signal "don't care" is put into the first byte of the fields not required.

## D. 3 Multi-page table <br> The MULTI-PAGE TABLE indicates all multi-pages currently transmitted with identification of the numbers of sub-pages used in each case. <br> D.3.1 Structure <br> The MULTI-PAGE TABLE is composed in accordance with the BASIC TOP TABLE by way of a coordinate system allocating a distinct position for each page of the 8 normal teletext magazines. <br> D.3.2 Coding <br> Coding of the MULTI-PAGE TABLE is in accordance with the code table given in [11]. The transmission codes are formed from the source codes by applying the Hamming $(8,4)$ code.

## D. 4 ADDITIONAL INFORMATION TABLE

For specific teletext pages, the ADDITIONAL INFORMATION TABLE contains additional indications, such as, for example, the title of a page and, in some cases, a so-called direct selection code.
D.4.1 Structure
The format of the page(s) containing the ADDITIONAL
INFORMATION TABLE(S) is adapted to that of the normal
teletext pages. An ADDITIONAL INFORMATION TABLE makes
use of 22 rows of a teletext page each having 40
columns. Each row may carry aditional information for
two individual teletext pages. In summary, from row 1
to $22,22 x 2=44$ additional items of information can
be indicated.
D. 4.2 coding

The information contained in the ADDITIONAL INFORMATION TABLE is coded according to [11].

For bytes 0 to 7 and 20 to 27, the transmission code is formed from the source code by applying the Hamming $(8,4)$ code; for bytes 8 to 19 and 28 to 39 , the transmission code is obtained by adding an odd parity bit to the 7-bit source code.

If fewer than 22 rows of the page containing the ADDITIONAL INFORMATION TABLE are used, the transmission of the unused rows may be dispensed with. In cases where the second field in the last row is not used, the "end" sign has to be entered. Otherwise, the lst field of the next following row must be marked with "don't care" and the second field with "end". Any rows still following the "end" sign are invalid.

Titles are entered at the left-hand edge of the title field, any remaining unneeded parts of the field are filled up with the sign "space" (code: 2/0). In titles, the sign "slash" (code: $2 / 15$ ) should not be used.

A maximum of two different pages may contain the same direct-selection code; the repetition of the same page numbers is not permissible.

The use of bytes 3 to 6 and 23 to 26 in rows 1 to 22 is reserved and not yet defined.

## E. TRANSMISSION

At the sending end the TOP tables are generated in accordance with the details specified in sections C. 1 to C.4.


#### Abstract

When inserting a teletext page into the transmission cycle, the teletext system must first insert that page and, susequentiy, update the TOP tables. When removing a page from the cycle, the teletext system must act inreverse: first up-dating the TOP tables and thereafter removing the corresponding pages. This prevents a decoder from awaiting the arrival of a page listed in the TOP tables after this page is no longer included in the cycle. The modification or up-dating of the BASIC TOP TABLE or one of the other TOP tables mentioned in the PAGE-LINKING TABLE is signalled to the decoding side by changing the numerical value in the sub-code area (bits b4, $\mathrm{b} 5, \mathrm{~b} 6$ ) in the page-header of the BASIC TOP TABLE (see c.1.2).

The BASIC TOP TABLE or, at least, the header of the page carrying that table should be refreshed at suitable time intervals, in order to convey any changes and up-dates of the TOP tables.


E. 1 Serial mode

## E.1.1 Basic TOP table

One BASIC TOP TABLE is transmitted for the complete teletext service. All displayable pages (having page numbers from 100 to 899) are marked in the BASIC TOP TABLE according to their current status. Rows 1 to 21 or 1 to 22 are transmitted completely; the page number for the transmission of the BASIC TOP TABLE is 1FO 3Fxx.

## E.1.2 Multi-page table

For the complete teletext service, only one MULTI-PAGE TABLE may be included in the transmission cycle.

If a MULTI-PAGE TABLE is transmitted in the transmission cycle, the page number (including the sub-code) and the identification of that table must be entered in the PAGE-LINKING TABLE. All multi-pages having decimal page numbers (100...899) are marked in the MULTI-PAGE TABLE in accordance with the number of associated sub-pages. Rows 1 to 20 are transmitted in their entirety.

## E.1.3 Additional information table

Additional information for all pages included in the transmission cycle is transmitted by use of a ADDITIONAL INFORMATION TABLE Or, if necessary, by use of several of these tables.

The page number (s) (including the sub-code) under which the ADDITIONAL INFORMATION TABLE(S) are transmitted have to be taken from the PAGE-LINKING TABLE.

## E. 2 Parallel mode

## E.2.1 Basic TOP table

For each teletext magazine an independent BASIC TOP TABLE has to be included in the transmission cycle.

Pages currently included in the magazine in question are marked in its BASIC TOP TABLE. Any rows of the BASIC TOP TABLE that are not relevant for the magazine in question, need not be transmitted.

The individual BASIC TOP TABLES are transmitted under page number YFO 3F7xx where $Y$ represents the magazine number concerned.

## E.2.2 Multi-page table

For each magazine included in the teletext transmission cycle, an independent MULTI-PAGE TABLE may be transmitted in the same or in another magazine.

Also in the "Parallel mode", a MULTI-PAGE TABLE can contain multi-page identifications for several magazines included in the transmission cycle.

If an allocated MULTI-PAGE TABLE is transmitted in connection with an individual magazine, the page number (including the sub-code) together with the identification of that table must be included in the PAGE-LINKING TABLE transmitted in the magazine in question.

All multi-pages with decimal page numbers of the allocated magazine are marked in the MULTI-PAGE TABLE in accordance with the number of associated sub-pages.

Any rows of the MULTI-PAGE TABLE that are not relevant for the allocated magazine(s) need not be transmitted.

## E. 2.3

Additional information table
For each magazine, independent additional information may be transmitted by use of an ADDITIONAL INFORMATION TABLE or, if necessary, by use of several tables in the same or in another magazine.

In the "Parallel mode", the ADDITIONAL INFORMATION TABLE can also contain additional information for several magazines transmitted in the the transmission cycle.

If additional information for a magazine is transmitted on one (or possibly several) ADDITIONAL INFORMATION TABLE(S), the page number(s) (including the sub-code) together with the identification(s) of these table(s) must be included in the PAGE-LINKING TABLE transmitted in the magazine in question.

## F. RECEIVER REQUIREMENTS

F. 1 Basic TOP table
Renewed read-in of BASIC TOP TABLE is necessary if the reception of that table is incomplete or erroneous.
In the case of the "parallel" magazine transmission, only those rows and columns of the BASIC TOP TABLE that carry relevant information for each individually transmitted magazine may be utilized.

## F. 2 Page-linking table

The bytes following a "don't care" sign or an "end" sign in the same field of the PAGE-LINKING TABLE are of no significance for the receiver.
The PAGE-LINKING TABLE has to be read in anew after every recognition of a sub-code change signalled in the page-header of the BASIC TOP TABLE. Renewed read-in is also necessary in the event of incomplete or erroneous reception.

## F. 3 Multi-page table

If the receiver detects for a specific page a multi-page identification, this is only valid if the page is simultaneously identified as a multi-page in the BASIC TOP TABLE.
The MULTI-PAGE TABLE has to be read in anew after every recognition of a sub-code change signalled in the page-header of the BASIC TOP TABLE. Renewed read-in is also necessary in the event of incomplete or erroneous reception.

In the case of "parallel" magazine transmission, only those rows and columns of the MULTI-PAGE TABLE that carry relevant information for the allocated magazine(s) may be utilized

## F. 4 Additional information table


#### Abstract

If the receiver detects a page carrying an ADDITIONAL INFORMATION TABLE the information is only valid for the page in question if the page is simultaneously marked in the BASIC TOP TABLE as a page with additional information.


If the receiver detects several currently valid pages identified by the same direct-selection code, that page with the smallest page number should be allocated to the directselection code.

The block title together with the titles of the associated groups and the titles of the pages identified by a directselection code form a unit which should be represented together for user information.

Renewed read-in is required in the event of faulty reception of the ADDITIONAL INFORMATION TABLE; until error-free reception, the page numbers in question could be displayed in place of the text carrying the user guidance (commentary row, title pages).

The ADDITIONAL INFORMATION TABLE has to be read in anew on every recognition of a sub-code change signalled by the pageheader of the BASIC TOP TABLE; renewed read-in is also required in the case of incomplete or erroneous reception.

The bytes following a "don't care" sign or an "end" sign in the same field of the ADDITIONAL INFORMATION TABLE are of no significance for the receiver.

## APPENDII 5

## RECOMMENDED CODE OF PRACTICE

FOR USE OF PAGE AND SUB-CODE ADDRESSES
There is a need for page headers to be available without any following data for time filling and page terminating applications. There is also a need for a page address to be available for internal use in decoders that will not occur in transmission. The following rules are recommended.

1. Any page address up to and including hexadecimal FE with a sub-code up to and including 3F7E, can be used for a page carrying data and can be specified as a linked page.
2. Pages address number $F F$ in any magazine, with any sub-code in the range 0000 to $3 F 7 E$ may be used for time filling and as a terminator. Pages with these addresses will not carry data for acquisition by a decoder.
3. Page header packets with the page address $F F$, sub-code $3 F 7 F$, will not be transmitted. This address will occur as a null link in packets with $\mathrm{Y}=27$ and $8 / 30$.

## APPENDIX 6

PAGE HEADER PACKET CONTROL BITS
C4 ERASE PAGE AND C8 UPDATE INDICATOR

The use of these control bits is subject to editorial judgement. The following is the expected effect.

1. When a page is selected for immediate display, that complete page is placed in memory.
2. Subsequent transmission of a page will overwrite the initially stored page. Any packets not overwritten will remain in memory for display.
3. Although ideally it would be desirable for all stored pages to be maintained updated, where a page is stored for subsequent rather than immediate display, the overwriting may only occur when a page is displayed.
4. The Update bit CB is used by the editor specifically to indicate that an update has occurred. The expected effect in the decoder is that, where a page display has been cancelled by an appropriate user key, ("cancel page", "picture" etc.), the setting of the Update bit will cause a prompt, which may involve automatic redisplay of the page. An application where this is standard practice is for News Flashes. The unnecessary or inappropriate setting of the Update bit can cause annoying redisplay of a page or News Flash that a user wishes to cancel. The setting of the Update bit is thus an editorial decision.
5. When the Erase bit C4 is set, that page is expected to be erased completely and a new page written into memory.
6. The effect of both of these bits applies to packets 1 to 28 inclusively.

## INDEPENDENT DATA SERVICES (DATAVIDEO FORMAT)

This protocol is based closely upon the specification given in Section 23 but with the variations indicated below:

Data channel addressing
Change only the numbers of allocated data channels:

| 5 | 1010 |
| :--- | :---: |
| 6 | 0110 |
| 13 | 1011 |
| 14 | 0111 |

S 23.4 Packet Address
Bytes 6, 7, 8, 4 bits data plus 4 bits Hamming protection.

S 23.5 Control bytes
Bytes 9 (CI) and 10 (CO), 4 bits data plus Hamming protection.

S 23.5.1 Packet continuity indicator
It represents a 6-bit number which is incremented modulo-64 with each new packet of the same address on the same data channel. It is not incremented on repeated transmissions of the same packet. Repeated packets may be interleaved with newly transmitted ones; in order to distinguish such repeated packets from new packets in case there is a discontinuity due to packet-losses, the interleaving pace must not be greater than 32.

Packet continuity index is transmitted in byte 9 and the firsi two data bits of byte 10; the least significant bit is transmitted first.

## S 23.5.2 Masking indicator

Bit $M$, which is the third data bit of byte 10 , indicates the presence of "masking": i.e.. when it has value 1 , this means that bytes 11 to 45 (useful data + CRC) are EX-ORed with the sequence (in hexadecimal notation):

AF AA 81 4A F2 EE 07 3A $4 F$ 5D 448670 9D 8343 BC 3F
EO F7 C5 CC 8253 B 479 F3 62 A4 71 B5 71311008.
It may be common practice to transmit the same packets once masked and once unmasked: this removes the adverse effects of critical bit combinations. Of course, masked and unmasked packets, when the information content is the same, are transmitted with the $s$ ame continuity index.

## S 23.5.3 Packet type indicator

Bit $L$, which is the fourth data bit of byte 10 , defines the packet type. It is used by higher levels of the protocol ar therefore it may be defined in a different way, according the type of protocol used. In one version of the protocol, this bit, when it is set to 1 , indicates the presence, in tr user data group, of a "packet length indicator". This is contained in byte 11, and indicates in binary code (least significant bit transmitted first, the last three bits reserved for future use) the number of following useful byte in the data body. In another version of the protocol, Bit I allows to distinguish between packets containing useful data ( $L=0$ ) and "control packets", i.e. packets containing service information such as access control information. The structu of auch packets is defined by higher levels of the protocol, and it is not part of the present specification.

## S 23.6 User data group

The remaining data bytes in the line, except the last two, (byte $11 t_{1}$ 43) constitute the data carried for users of the service, bearing tha Packet Address in that Data Channel. There are 33 data bytes available.

S 23.7 Cyclic Redundancy Check Word
The last two bytes contain a Cyclic Redundancy Check on the User Data group.

S 23.7.1 Check Word Generation
The data to be checked is considered as a polynomial in $x$ wi the highest degree term transmitted first and the term of degree zero last.

This is divided, using modulo-2 arithmetic, by the polynomia
$x 16+\times 12+x 5+1$
The remainder from this process, with the highest term transmitted first, is the CRC. Before transmission the bits are complemented (i.e. XORed with FF hex).

## S 23.7.2 Check result

The register of figure 24 is set to $0 s$. The serial data followed by the CRC is then entered. The check is satisfied if the register again contains all os.



Figure 1
television signal line numbers and data timing


| Page Number Units | Page Number Tens | Page SubCode | Page SubCode \& Control bit c4 | $\begin{aligned} & \text { Page Sub- } \\ & \text { Code } \end{aligned}$ | Page SubCode, C5 \& C6 | Control Bits $\mathrm{C} 7-\mathrm{C} 10$ | ```Control Bits C11 - C14``` | 32 Character Bytes Odd Parity |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PDPDPDPD | PDPDPDPD | PDPDPDPD | PDPDPDPC4 | PDPDPDPD | PDPDPC5PC6 | PC7PC8PC9PC10 | PC11PC12PC13PC14 | DDDDDDDP | DDDDDDDP | etc |

[^0]

Packet $\mathrm{x} / 26:$ Di-D6 Address, D7-11 Mode, D12-D18 Data, P1-P6 Hamming(24, 18) protection bits
Packets $\mathrm{X} / 28$ \& 29: 13 Three Byte Data Groups in each Packet $\mathrm{P}=\mathrm{Hamming}$ Bit $\mathrm{D}=\mathrm{Data}$ Bit

| X/27 | Designation Codes 0000 to 0011 <br> (for codes 0100 <br> to 0111 see <br> Fig. 3 parts 2\&3) | First Group of Six Bytes containing Relative Magazine number, page number and Page Sub-Code | Five further Groups of Six Bytes each plus byte 43 containing link control Data | When Designation Code is 0000, final two bytes are a basic page check digit |
| :---: | :---: | :---: | :---: | :---: |


| 8/30 | $\mathbf{P} \quad \mathrm{D} P \mathrm{D}$ P D P D | One Six-Byte Group Coded | Network Ident, Time Offset, Mod Julian Date, Co-Ordinated |
| :---: | :---: | :---: | :---: |
| Format | Designation Code | as $\mathrm{X} / 27$, Designation Codes | Universal Time, 18t and 2nd Short Programme Labels, |
| 1 | $\mathbf{X} 000$ | 0000 to 0011. See NOTE | bytes 26 to end of packet, Data for Direct Display |


| 8/30 | $\mathbf{P} \mathbf{D} \mathbf{P} \mathbf{D} \mathbf{P} \mathbf{D} \mathbf{P} \mathbf{D}$ | One Six-Byte Group Coded | Programme | Data for Direct Display |
| :---: | :---: | :---: | :---: | :---: |
| Format | Designation Code | as $\mathrm{X} / 27$, Designation Codes | Identification | Bytes 26 to end |
| 2 | X 1000 | 0000 to 0011. See NOTE | Data | of packet |

NOTE These six bytes have the same format as bytes 6 to 11 of the page header packet, see Figure 2 . The bits corresponding to the Control Bits $C 4, C 5$ and $C 6$ in the page header packet are used in this sequence to change the magazine number from that in bytes 4 and 5 of the packet $x / 27$. Setting any of these bits to 1 ' complements the corresponding magazine number bit.

$P=$ Hamming $(24,18)$ Protection Bit, $D=$ Data Bit

First Three Bytes of a six-Byte Group
Data Bits 1-11: Link Control Data (see Figure 3 part 3),
Data Bits 12-14: Relative Magazine Number (see NOTE),
Data Bits 15 - 18: Page Number Tens
Second Three Bytes of a six-Byte Group
Data Bit 1: Set to 0,
Data Bits 2 - 5: Page Number Units,
Data Bits 6 - 18: Page Sub-code

NOTE The three Relative Magazine bits change the magazine number from that in bytes 4 and 5 of the packet $\mathrm{X} / 27$. Setting any of these bits to 1 complements the corresponding magazine bit.

In all cases the LEAST SIGNIFICANT bit is transmitted first

PACKETS WITH Y=27 DESIGNATION CODES 0100 TO 0111


Figure 3 part 3 LINK CONTROL DATA FOR CONPOSITIONAL APPLICATIONS


Note 1 These bits are reserved.
Note 2 Each 4-bit number is incremented by one prior to transmission. The pairs of 4 -bit numbers are assembled into bytes and the bytes are transmitted least significant bit first.

Bytes 13 to the end of the packet....


## NOTE: Each byte contains 4 bits data plus 4 bits Hamming ( 8,4 ) protection.

The data bits in each data word are transmitted most significant bit first.
Groups of 4 data bits are assembled, with Hamming( 8,4 ) protection bits, into bytes.

$D=$ data bit $P=$ parity bit
 Allocation Flag except for Continuity Indicator where D1 to D14 represent a binary number incremented at each change in the page content.

Bytes 38 to 45 are reserved.
Packets with $Y=17$ to $Y=26$ are reserved.


First PDC packet X/26


```
Bits D1 - D6 Data word 'A'
Bits D7 - D11 Mode
Bits D12 - D18 Data Word 'B'
Bits P1 - P6 Hamming(24,18) Protection
```

Last PDC packet X/26

| Clock Run-in | Framing <br> Code | Magazine and <br> Packet Address | Design- <br> ation <br> Code | Last PDC Data Group | Terminating Data Group(s) and CRC |
| :--- | :--- | :--- | :--- | :--- | :--- |

## 16 bit SHIFT REGISTER



In the example shown a 16 bit shift register has as input the modulo-2 sum of an external input and the contents of the 7 th, $9 t h$, 12th and 16th stages of the register. Initially the register is cleared to "all zeros". During a sequence of 8192 clock pulses, the first 24 character bytes (192 bits) of the page header packet and the following character bytes of packets with Y up to 25 , in conventional transmission order, form the input. Any absent packets are considered to contain the character SPACE (2/0) throughout. In each byte, the bit order is bs to bl inclusive. This order, that is the reverse of that used in the transmission sequence, is to facilitate decoder operation where the data used is stored in the page memory.
At the transmitting end of the generating process the contents of the register are the basic page check word and it is transmitted along the register beginning with the bit held in the first stage.

The transmission order for the two byte group resulting from the 16 bit cyclic redundancy check on the page is bits 9 to 16 followed by bits 1 to 8 inclusive.

| MODE | FUNCTION AND DATA BITS 1 to 7 |
| :---: | :---: |
| 00000 | FULL SCREEN COLOUR |
|  | Bits 1 to 5 least significant bit first invoke an entry from the Colour Map. Bits 6 and 7 are set to 0 |
| 00001 | FULL ROW COLOUR |
|  | Bits 1 to 5 least significant bit first invoke an entry from the Colour Map. Bits 6 and 7 set to 0 single row; Bits 6 and |
|  | 7 set to 1 contiguous rows; Other Codes, ignore data bits USED BY SANSKRIT-RELATED LANGUAGES |
| 00011 | Further information is available in a supplementary specification [3] |
| 00100 | SET CURSOR POSITION |
|  | Cursor position, left to right is defined by codes $4 / 0$ to 6/7, other codes are to be interpreted as 'No Cursor' |
| 00101 | StART SCROLI REGION |
|  | Data bits 1 to 5 defines full row scroll region colour, |
| 00110 | END SCROLL REGION |
|  | Data bits 1 to 5 define full row colour from next row to lower border area inclusively, bits 6 and 7 set to 0 |
| 00111 | ADDRESSES DISPLAY ROW 0 |
|  | When data bits are set to O Section 14.6.4.1 applies |
| 01000 | PDC DATA |
|  | Row Address bits define Country, Data bits 1 to 6 define Source |
| 01001 | PDC DATA |
|  | Row Address bits define Date Months, Data bits 1 to 6 define Day |
| 01010 | PDC DATA |
|  | Row Address bits define Cursor Row, Data bits 1 to 6 define Starting Time Hours |
| 01011 | PDC DATA |
|  | Row Address bits define cursor Row, Data bits 1 to 6 define Finishing Time Hours |
| 01100 | PDC DATA |
|  | Row Address bits define Cursor Row, Data bits 1 to 6 define Local Time Offret |
| 01101 | PDC DATA |
|  | Row Address bits invoke Series Code |
|  | Function, Data bits 1 to 7 provide 'Series Code' |
| 01110 to | NOT ASSIGNED |
| 11110 | Ignore data bits |
| 11111 | TERMINATOR |
|  | For packets with Y $=26$ |

Figure 9
PACKETS WITH Y $=26$ ROW ADDRESS GROUP 40 to 63

MODE DESCRIPTION CODE

| MODE | FUNCTION AND DATA BITS 1 to 7 |
| :---: | :---: |
| 00000 | FOREGROUND COLOUR <br> Bits 1 to 5 least significant bit first, invoke an entry from the colour map. Bits 6 and 7 are set to 0 |
| 00001 | G1 SET Bits 1 to 7 define the character. Block Mosaic is default character set |
| 00010 | G3 SET Bits 1 to 7 define the character. Smoothed Mosaic is the default character set |
| 00011 | background colour <br> Bits 1 to 5 least significant bit first, invoke an entry from the colour map. Bits 6 and 7 are set to 0 |
| 00100 | LATCBING SHIFT <br> See Figure 10 part 2 |
| 00101 | SINGLS SHIFT |
|  | See Figure 10 part 3 |
| 00110 | PDC DATA <br> Address bits define Cursor <br> Position in Row, Data bits 1 to 7 see Section 20.4.5.2 |
| 00111 | ADDITIONAL FLASH CONTROLS See Section 14.6.8 |
| 01000 | GO SET MODIFIED DESIGNATION Bita 1 to 7 define a character set, see figure 14 |
| 01001 | G1 SET MODIFIED DESIGNATION <br> Bits 1 to 7 define a character set, see figure 14 |
| 01010 | G2 SET MODIFIED DESIGNATION Bita 1 to 7 define a character set, see figure 14 |
| 01011 | G3 SET MODIFIED DESIGNATION Bits 1 to 7 define a character set, see figure 14 |
| 01100 | NON-SPACING ATHRIBUTES |
|  | See Section 14.6 |
| 01101 | 1 st DRCS |
|  | Bits 1 to 7 define a character from the 1st DRCS |
| 01110 | 2nd DRCS 7 define a character from the 2nd DRCS |
| 01111 | Bits 1 to 7 define a character from the 2nd DRCS CHARACTERS FROM G2 SET |
|  | Bits 1 to 7 define a character from the G2 Supplementary Set |
| 10000 to | DIACRITICAL MARRS FROM COLUN 4 OF G2 SET |
| 11111 | Bits 1 to 7 define the associated GO Primary Set Character |


| DATA BITS 1 TO 7 | FUNCTION when space mode description bits are set to 00100 |
| :---: | :---: |
| 0000000 | Latching Shift to Invoked or Default G0 Character Set |
| 0000001 | Latching Shift to Invoked or Default Gl Character Set |
| 0000010 | Latching Shift to Invoked or Default G2 Character Set |
| 0000011 | Latching Shift to Invoked or Default G3 Character Set |
| 0000100 | Latching Shift to lst DRCS |
| 0000101 | Latching Shift to 2nd DRCS |
| $\begin{aligned} & 0000110 \text { to } \\ & 1111110 \end{aligned}$ | Invocation of and Latching to character sets defined in Figure 14 |
| 1111111 | Cancel Latching Shift or Invocation and Latching Shift |

The effect of a shift or combined invocation and shift is cancelled by transmission of a further shift or a combined invocation and shift.

Figure 10 part 2 PACKETS WITH Y $=26$
LATCHING SHIFTS AND COMBINED INVOCATION AND LATCHING SHIFTS

| DATA BITS 1 to 7 | FUNCTION when space group mode description bits are set to 00101 |
| :---: | :---: |
| 0000000 | Single Shift to Invoked or Default GO character set |
| 0000001 | Single Shift to Invoked or Default Gl character set |
| 0000010 | Single Shift to Invoked or Default G2 character set |
| 0000011 | Single Shift to Invoked or Default G3 character set |
| 0000100 | Single Shift to lst DRCS |
| 0000101 | Single shift to 2nd DRCS |
| $\begin{aligned} & 0000110 \text { to } \\ & 1111110 \end{aligned}$ | Invocation of and Single Shift to character gets defined in Figure 14 |
| 1111111 | Reserved |

The effect of a shift or combined invocation and shift is cancelled by transmission of a further shift or a combined invocation and shift.

```
Figure 10 part 3 PACRETS WITH Y=26
    SINGLE SHIFTS AND COMBINED INVOCATION AND SINGLE SHIFT
```

| 18 | 17 | 16 | 15 | Bit Numbers |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 | 8 bit bytes, 7 bits data and 1 parity bit |
| 0 | 0 | 0 | 1 | 8 bit bytes, with 8 data bits each |
| 0 | 0 | 1 | 0 | 3 groups of 8 bit bytes, 18 data bits, 6 Hamming bits |
| 0 | 0 | 1 | 1 | 8 bit bytes, 4 data bits, 4 Hamming bits |
| The interpretation of other combinations of bit values |  |  |  |  |
| is reserved |  |  |  |  |

Figure 11 part 1 Interpretation of the first group of 18 data bits in a packet with $Y=28$ of a page with designation code 0000 or 0010

| 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | Bit Numbers |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 |  | 0 | 0 | 0 | 0 | 0 | 0 | Basic Page with Standard Character Position and Row Format |
| 0 | 0 | 1 | * | * | * | * | * | Overwriting Paeudo Page |
| 0 | 1 | 0 | * | * | * | * | * | Scrolling Pseudo Page |
| 0 | * | * | 0 | 0 | 0 | 0 | 0 | Page with Standard Character Position and Row Format |
| 0 | * | * | 1 | * | * | * | * | Page Format Extension, Not Last of Group |
| 0 | * | * | 0 | * | * | * | 0 | Page Format Extension, Last of Group |
| 0 | * | * | * | * | * | 0 | 0 | Page Format Extension, Left Hand Page |
| 0 | * | * | * | * | * | 0 | 1 | Page Format Extension, 2nd from Left Page |
| 0 | * | * | * | * | * | 1 | 0 | Page Format Extension, 3rd from Left Page |
| 0 | * | * | * | - | * | 1 | 1 | Page Format Extension, 4th from Left Page |
| 0 | * | * | * | 0 | 0 | * | * | Page Format Extension, Top Row of Pages |
| 0 | * | * | * | 0 | 1 | * | * | Page Format Extension, 2nd from Top Row |
| 0 | * | * | * | 1 | 0 | * | * | Page Format Extension, 3rd from Top Row |
| 0 | * | * | * | 1 | 1 | * | * | Page Format Extension, 4th from Top Row |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | DRCS Downloading, First Group |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | DRCS Downloading, Second Group |
| 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | Idiographic Character Downloading, lst Group |
| 1 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | n $n$ n 2nd Group [4] |
| 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | Reformatted Data |
| 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | Terminal Equipment Addressing Page |
| 1 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | Musical Sound Data [5] |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | Geometric Data Profile '0' Option 1 |
| 1 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | Geometric Data Profile '0' Option 2 |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | Geometric Data Profile '1' Option 1 |
| 1 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | Geometric Data Profile '1' Option 2 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | Geometric Data Profile '2' Option 1 |
| 1 | 0 | 0 | 0 | 1 | 1 | 1 | 0 | Geometric Data Profile '2' Option 2 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | Geometric Data Profile '3' Option 1 |
| 1 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | Geometric Data Profile '3' Option 2 |
| 1 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | Photographic Data Option 1 |
| 1 | 0 |  | 1 | 1 | 0 | 0 | 0 | Photographic Data Option 2 |
|  | = |  | C | re |  | lue |  | The interpretation of other combinations of bit values is reserved |

Figure 11 part 2 Interpretation of the first group of 18 data bits in a packet with $Y=28$ of page with designation code 0000 or 0010

This interpretation is only assigned to the values of bits 1 to 8 , when bits 9 to 14 are set to 0 .


The downloading bit combinations are transmitted least significant bit first in the order red, green, blue. Reduced Intensity Yellow, as an example, is transmitted as 111011100000.

CHARACTER MODE 1 (basic) Mode Identification: 0000; 12 cols. $x 10$ rows $x 1$ 20 D-byte Group of a Packet [A,B,C,D....T]: One PTU

| Column bits $12-7$ | $6-1$ | Rows | Dot positions defined as 1 will be |
| :---: | :---: | :---: | :---: | :--- |
| A | B | 1 | displayed in Foreground Colour. |
| C | D | 2 | Dot positions defined as 0 will be |
| E | F | 3 | displayed in Background Colour. |
| $:$ | $:$ | $:$ |  |
| S | T | 10 |  |

CHARACTER MODE 2 Mode Identification: 0001; $12 \times 10 \times 2$
2 PTUs for a character
20 D-byte Group of Packet $N$ [A1,B1,C1,D1....Tl]: One PTU

| Column bits | $12-7$ A1 | $6-1$ B2 | Rows 1 | Dots are defined by a bit in Character Code "X1" and an equivalent bit from |
| :---: | :---: | :---: | :---: | :---: |
| Character | C1 | D2 | 2 | Character Code "X2". |
| Code "X1" | E1 | F2 | 3 |  |
|  | $\begin{aligned} & \text { S1 } \end{aligned}$ | $\begin{aligned} & : \\ & \mathbf{T} 2 \end{aligned}$ | $10$ | The bit from "X1" specifies the least aignificant bit from the DRCS Colour LookUp Table (DCLUT) |

Next 20 D-byte Group of Packet $N$ [A2, B2,C2,D2....T2]: One PTU

| Column bits | $12-7$ | $6-1$ | Rows | The bit from "X2" specifies the most |  |
| :--- | ---: | ---: | :---: | :--- | :--- |
|  | A2 | B1 | 1 | gignificant bit of the DRCS Colour Look- |  |
| Character | C2 | D1 | 2 | Up Table (DCLUT). |  |
| Code "X2" | E2 | F1 | 3 |  |  |
|  | $\vdots$ | $\vdots$ | $\vdots$ |  |  |
|  | S2 | T1 | 10 |  |  |

CHARACTER MODE 3 Mode Identification: 0101; Two gets of $6 \times 10 \times 1$ 20 D-byte Group of a Packet $N \quad[A, X, B, X, C, X \ldots J, X]$ : one PTU

| Column | 1st DRCS | 2nd | DRCS | Dots positions defined as 1 will be displayed in Foreground Colour. Dots positions defined as 0 will be displayed in Background Colour. |
| :---: | :---: | :---: | :---: | :---: |
|  | bits 12-7 | 6-1 | Rows |  |
|  | A B | - | 1 |  |
|  | C | - | 3 |  |
|  | J |  | $1{ }^{\text {: }}$ |  |

20 D-byte Group of a Packet $N \quad[X, A, X, B, X, C, \ldots, X, J]:$ One PTU



CHARACTER MODE 5 Mode Identification: 0111; One set $6 \times 10 \times 4$ Two PTUs required for a character


| Column bits <br> Character Code X1 | 12-7 | 6-1 | Rows | The combination of bits 4,3,2,1 is used as |
| :---: | :---: | :---: | :---: | :---: |
|  | A3 | A4 | 1 | the binary address to colour tables 3 and |
|  | B3 | B4 | 2 | 4 of the colour map, figure 12. |
|  | C3 | C4 | 3 |  |
|  | : ${ }_{\text {J }}$ | $\stackrel{1}{5}$ | 10 |  |



Next 20 D-byte Group of a Packet $N\{X, A 1, X, A 2, X, B 1, X, B 2, X, C 3, \ldots, X, E 2$,$\} :$ One PTU

| Column bits 1 | 12-7 | $6-1$ A1 | Rows <br> 1 | Dots are defined by equivalent pair 1 and 2. |
| :---: | :---: | :---: | :---: | :---: |
| Character of | X | A2 | 2 | 1 specifies the least significant bit from |
| the Second | X | B1 | 3 | the DRCS Colour Look-up Table. |
| Set | X | E2 | 10 | 2 specifies the most significant bit from |
|  | X | E2 | 10 | the DRCS Colour Look-up Table. |

CHARACTER MODE 7 Mode Identification: 1000; One set $6 \times 5 \times 4$
20 D-byte Group of a Packet $N$ [A1,A2,A3,A4,B1,B2....E3,E4]: One PTU

| Column bits | $12-7$ | $6-1$ | Rows | The combination of bits $4,3,2,1$ is used as |  |
| :--- | ---: | ---: | :---: | :--- | :--- |
|  | A1 | A2 | 1 | the binary address to colour tables 3 and |  |
| Character | A3 | A4 | 2 | 4 of the colour map, figure 12. |  |
| Code X1 | B1 | B2 | 3 |  |  |
|  | $:$ | $\vdots$ | 4 |  |  |
|  | E1 | E2 | 10 |  |  |

Figure 13 part 3 DATA FORMAT FOR DRCS DOWNLOADING

| SET | CODE | NAME AND CODE TABLE FIGURE NUMBER |
| :---: | :---: | :---: |
| $\begin{aligned} & 0 \text { to } \\ & 5 \end{aligned}$ | $\begin{aligned} & 0000000 \text { to } \\ & 0000101 \end{aligned}$ | Reserved for other applications |
| 6 | 0000110 | 1st Latin Primary Set 15\&17 |
| 7 | 0000111 | lst Latin Supplementary Set 16 |
| 8 | 0001000 | Block Mosaic Set 18 |
| 9 | 0001001 | Smoothed Mosaic Set 19 |
| 10 | 0001010 | Arabic Primary Set 21 |
| 11 | 0001011 | Arabic Supplementary Set 21 |
| 12 | 0001100 | Cyrillic Primary Set 23,25\&26 |
| 13 | 0001101 | Cyrillic Supplementary Set 24 |
| 14 | 0001110 | Reserved |
| 15 | 0001111 | Reserved |
| 16 | 0010000 | Greek Primary Set 27 |
| 17 | 0010001 | Greek Supplementary Set 28 |
| 18 | 0010010 | Reserved |
| 19 | 0010011 | Reserved |
| 20 | 0010100 | Hebrew Primary Set 22 |
| 21 | 0010101 | Hebrew Supplementary Set |
| 22 to | 0010110 to | Reserved |
| 37 | 0100101 |  |
| 38 | 0100110 | 2nd Latin Primary Set 15817 |
| 39 | 0100111 | 2nd Latin Supplementary set |
| 40 | 0101000 | 3rd Latin Primary Set 15@17 |
| 41 | 0101001 | 3rd Latin Supplementary Set |
| 42 | 0101010 | Cyrillic and Latin Primary sets 26 |
| $43 \text { to }$ $54$ | $\begin{aligned} & 0101011 \text { to } \\ & 0110110 \end{aligned}$ | Reserved |
| 55 | 0110111 | 4th Latin Primary Set 15\&17 |
| 56 | 0111000 | 4th Latin Supplementary Set |
| 57 | 0111001 to | Reserved |
| 69 | 1000101 | Yugoslav (\& other former states) |
| 70 | 1000110 | " Latin Primary Set |
| 71 | 1000111 | " Latin Supplementary Set |
| 72 to | 1001000 to | Reserved for other applications |
| 75 | 1001011 | Yugoslav (\& other former states) |
| 76 | 1001100 | * Cyrillic Primary Set |
| 77 | 1001101 | " Cyrillic Supplementary Set |
| $\begin{aligned} & 78 \text { to } \\ & 127 \end{aligned}$ | $\begin{aligned} & 1001110 \text { to } \\ & 1111111 \end{aligned}$ | Reserved for other applications |

Entries identified as "Reserved for other applications" may not designate a character set. Entries identified as "Reserved" have been allocated to other character sets or writing systems.

Note: Set 70 is identical to Set 6 and Set 71 is identical to Set 7


Notes:
Optional variations apply to marked positions (see Figure 17).
The character in position $2 / 0$ is to be interpreted as SPACE.
The character in position $7 / 15$ occupies an area equivalent to that of any character which does not include a descender. It is thus a rectangle surrounded by the background colour.


Notes:
Column 4 contains diacritical marks for association with GO set character
The character in position $7 / 15$ occupies an area equivalent to that of any character which does not include a descender. It is thus a rectangle surrounded by the background colour.

The character in position $2 / 0$ is to be interpreted as SPACE.
The characters in positions 5/6, 5/7 and 5/8 are in addition to those specified in ISO 6937/2.

The character in position $5 / 6$ represents the European Currency symbol.
Blank positions are "transparent" permitting the display of any character otherwise overwritten.

|  | $\begin{aligned} & \text { TABLE } \\ & \text { POSITION } \end{aligned}$ | CZECH/ slovak | english | FRENCH belgian | GERMAN | Italian | POLISH |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 2/3 | * | £ | E | \# | E | * |
| 2 | 2/4 | $\dot{\text { u }}$ | \$ | I | \$ | \$ | ¢ |
| 3 | 4/0 | c | e | a | S | E | 9 |
| 4 | 5/11 | t' | $\leftarrow$ | ë | A | - | $\dot{\mathbf{z}}$ |
| $\because$ | 5/12 | z | $\frac{1}{2}$ | ¢ | 0 | ¢ | $s$ |
| 6 | 5/13 | $\underline{y}$ | $\rightarrow$ | ù | 3 | $\rightarrow$ | も |
| - | 5/14 | 1 | $\uparrow$ | 1 | $\wedge$ | $\uparrow$ | $\varepsilon$ |
| 8 | 5/15 | ř | * | * | - | * | $\delta$ |
| 9 | 6/0 | E | - | è | - | ù | e |
| - 0 | 7/11 | a | $\frac{1}{4}$ | a | a | à | $\dot{z}$ |
| 1. | 7/12 | と | 1 | $\delta$ | \% | ò | 6 |
| 12 | 7/13 | ú | 3 | û | ü | è | 1 |
| 13 | 7/14 | \% | + | ¢ | B | i | $\Sigma$ |


| $\begin{aligned} & \text { TABLE } \\ & \text { POSITION } \end{aligned}$ | PORTUGUESE SPANISH | RUMANIAN | SERBIAN CROATIAN SLOVENIAN | SWEDISH <br> FINNISH | TURKISH |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 2/3 | ¢ | \# | \# | * | TL |
| 2/4 | \$ | $\square$ | \$ (see | $\square$ | $\overline{\mathbf{g}}$ |
| 4/0 | 1 | T |  | $E$ | $\dot{I}$ |
| 5/11 | a | $\AA$ | ć | Ä | Ş |
| 5/12 | E | Ş | z | Ö | Ö |
| 5/13 | i | Ă | Đ | $\AA$ | Ç |
| 5/14 | $\sigma$ | I | S | ¢ | U |
| 5/15 | u | 1 | ë | - | $\overline{\mathbf{G}}$ |
| 6/0 | $\varepsilon$ | ¢ | c | E | 1 |
| 7/11 | ü | a | $\varepsilon$ | a | 8 |
| 7/12 | ก | \% | ž | 0 | \% |
| 7/13 | è | a | d | 3 | ¢ |
| 7/14 | a | 1 | \% | ü | $\ddot{\text { u }}$ |

The character TL is the Turkish currency sign
Note: $\ddot{E}$ has been implemented in some intregrated circuits

| $\begin{aligned} & \text { OPTION NUMBER } \\ & \text { AND } \\ & \text { C12 C13 C14 } \end{aligned}$ |  | CHARACTER SET NUMBER AND <br> ASSOCIATED ALPHABETS |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 6 | 38 | 55 | 70 |
| 1 | 000 | English | Polish | English |  |
| 2 | 001 | German | German | German | Albanian |
| 3 | 010 | Swedish Finnish | Swedish Finnish | Swedish Finnish | Hungarian |
| 4 | 011 | Italian | Italian | Italian |  |
| 5 | 100 | French | French | French | Slovakian |
| 6 | 101 | Portugese Spanish |  | Portugese Spanish |  |
| 7 | 110 | Czech Slovak | Czech Slovak | Turkish | Serbian Croatian Slovenian |
| 8 | 111 | Reserved | Reserved | Rese ved | Rumanian |

See Figure 14 Character Set Numbers for codes in packets with $Y=28$

Figure 17 part 2 IATIN ALPPABET GO CHARACTER SET OPTIONS




Notes:
Columns 4 and 5, indicated by (1), contain the corresponding characters from the GO set.

Shaded areas of characters in columns 2, 3, 6 and 7 are displayed in the foreground colour.

Bit allocations are the same as for the Primary character set, see Figure 15 .

Figure 18 MOSAIC GRAPHICS G1 CEARACTER SET (shown in contiguous form)


Notes:
The bit allocations are the same as for the Primary character set, see Figure 15.

The character at position $2 / 15$ represents the activation of alternate dots 0 : the matrix.
(1) This character permits display of Background colour.
(2) No character is assigned to these codes. Should they occur, the Level 1 character shall be displayed.

|  | 0 | 1 |
| :---: | :---: | :---: |
| 0 | Alpha Black (4) | Mosaic Black (4) |
| 1 | Alpha Red | Mosaic Red |
| 2 | Alpha Green | Mosaic Green |
| 3 | Alpha Yellow | Mosaic Yellow |
| 4 | Alpha Blue | Mosaic Blue |
| 5 | Alpha Magenta | Mosaic Magenta |
| 6 | Alpha Cyan | Mosaic Cyan |
| 7 | Alpha White (1) | Mosaic White |
| 8 | Flash | Conceal (2) |
| 9 | Steady (1,2) | Contiguous Mosaic (1,2) |
| 10 | End Box (1,3) | Separated Mosaic (2) |
| 11 | Start Box (3) | (5) |
| 12 | Normal Height (1,2) | Black Background (1,2) |
| 13 | Double Height | New Background (2) |
| 14 | Double Width (4) | Hold Mosaic (2) |
| 15 | Double size (4) | Release Mosaic (1) |

## Notes:

(1) Presumed at the start of each display row
(2) Action "Set At", others are "Set After"
(3) Two consecutive codes transmitted, action takes place between them
(4) No action at level 1
(5) Used in connection with Data for Processing Associated with Displayable Pages (see Section 19.2), and in some countries for language switching.

|  | 0 | 2 | O | - | 5 |  |  | 1916 | 51 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 5 | U | 1 | N | $\bigcirc$ | ! |  |  | 0111 | 71 |
| n | ш | n | W | $\div$ | $\stackrel{\square}{-}$ |  |  | 1016 | El |
| O | I | ə | T | K | $\therefore$ |  |  | 10112 | Z1 |
| ¢ | Y | ? | X | 3 | $\stackrel{\square}{7}$ |  |  | 11011 | 11 |
| Z | ! | Z | I | ? | 2 |  |  | 0101 | 01 |
| $\Lambda$ | ! | X | I | 6 | $\sim$ |  |  | 1001 | 6 |
| X | 4 | X | H | $V$ | - |  |  | 0000 | 8 |
| M | 8 | M | $\bigcirc$ | $\wedge$ | ร. |  |  | 110 | L |
| $\Lambda$ | J | $\Lambda$ | H | 1 | 5 |  |  | 01610 | 9 |
| n | $\partial$ | П | 3 | 0 | $i$ |  |  | 010 | 5 |
| 7 | p | L | C | 3 | $!$ |  |  | 01010 | 7 |
| S | 0 | S | 0 | d | $!$ |  |  | 1100 | $E$ |
| I | 9 | y | g | $\lambda$ | 1 |  |  | 0100 | 2 |
| b | E | O | V | 1 | 3 |  |  | 1000 | 1 |
| d | ? | d | - | - |  |  |  | ${ }^{2} 000$ | 0 |
|  |  |  |  |  |  |  |  | - ${ }^{\text {ama }}$ |  |
|  | 0 | 1 | 0 | - | 0 | 1 | 0 | 99 |  |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 59 |  |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | L9 |  |
| $L$ | 9 | 5 | 7 | E | 2 | 1 | 0 |  |  |


| - | - | \# | $\bigcirc$ | j | 1 |  |  | 11 |  | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | $?$ | ب- | $\sim$ | $>$ |  | N |  | 01 | 11 | 7 |
| ? | 2 | $\sim$ | $\sim$ | $=$ | - | - | 1 | 101 | 1. | 1 E |
| $-$ | 2 | - | - | $<$ | , |  |  | 00 | 11 | Z |
| $\downarrow$ | $?$ | $\because$ | ! | ; | $+$ |  |  |  | 01 | 1 |
| $\sim$ | 「 | 5 | 5 | : | * |  |  | 0110 | ${ }^{2}$ | 0 |
| $\Gamma$ | $\Omega$ | 5 | $\bigcirc$ | 6 | ) |  |  | 100 | 01 | 6 |
| 7 | 6 | 9 | $\stackrel{\square}{\circ}$ | 8 | ( |  |  | 00 | 0 | 8 |
| $\sqrt{5}$ | $\sigma$ | 9 | 1 | $L$ | ร |  |  | 111 |  | $L$ |
| $\bigcirc$ | ? | $\sigma$ | 7 | 9 | $\cdots$ |  |  | 011 | 10 | 9 |
| T | $\sim$ | $\sigma$ | $\bigcirc$ | 5 | $\%$ |  |  | 101 | 10 | 5 |
| $\square$ | 5 | - | $\div$ | 7 | \$ |  |  | 001 | 10 | 7 |
| - | 5 | - | $\dot{\sim}$ | $\varepsilon$ | 7 |  |  |  | 00 | $E$ |
| - | $\underline{5}$ | ! | - | Z | " |  |  |  | 00 | 2 |
| 2 | ! | $r$ | 3 | 1 | i |  |  | 100 | 00 | 1 |
| $s$ | - | ¢ | $\stackrel{\square}{?}$ | 0 | dS |  | Yo |  | 0 | 0 |
|  |  |  |  |  |  |  |  | $\square$ | \% ${ }^{5}$ |  |
| 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 99 |  |  |
| T | 1 | 0 | 0 | 1 | 1 | 0 | 0 |  |  |  |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | L9 |  |  |
| $L$ | 9 | 5 | 7 | $E$ | 2 | 1 | 0 |  |  |  |


|  | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 |  | 0 | e | P | $\stackrel{ }{*}$ | 1 |
| 1 | $!$ | 1 | A | Q | J | 0 |
| 2 | " | 2 | B | R | 1 | I |
| 3 | £ | 3 | C | S | 7 | $\eta$ |
| 4 | \$ | 4 | D | T | 1 | 9 |
| 5 | $\%$ | 5 | E | U | 1 | $\gamma$ |
| 6 | \& | 6 | F | v | 1 | צ |
| 7 | , | 7 | G | W | $\pi$ | $p$ |
| 8 | $($ | 8 | H | X | $\varphi$ | 7 |
| 9 | ) | 9 | I | Y | ' | $\emptyset$ |
| 10 | * | : | J | 2 | 7 | $\Pi$ |
| 11 | + | ; | K | $\leftarrow$ | J | (a) |
| 12 | , | < | L | $\frac{1}{2}$ | 7 | IV |
| 13 | - | = | M | $\rightarrow$ | $\square$ | $\frac{3}{4}$ |
| 14 | - | $>$ | N | $\dagger$ | D | $\div$ |
| 15 | 1 | ? | 0 | \# | 1 | $\square$ |

Notes:
The character (a) in position $7 / 11$ is the Shekel currency sign Bits 1 to 4 define rows, bits 5 to 7 define columns.


Notes:
The character in position $2 / 0$ is to be interpreted as SPACE Characters marked (1) are National variations, see Figure 2!


Notes:
Column 4 contains diacritical marks for association with another character addressed to the same presentation position.

The character in position $2 / 0$ is to be interpreted as SPACE.
The character in position $5 / 6$ represents the European Currency symbol.
Blank positions are "transparent", permitting the display of any character otherwise overwritten.

| CODE <br> TABLE <br> POSITION | SERBIAN CROATIAN （Option 1） | RUSSIAN BULGARIAN （Option 2） |
| :---: | :---: | :---: |
| $2 / 6$ | \＆ | H |
| 4／0 | 4 | 10 |
| 4／10 | J | Й |
| 5／1 | K | G |
| 5／6 | B | I |
| 5／7 | $\Gamma$ | B |
| 5／8 | b | b |
| 5／9 | 历 | ＇b |
| 5／11 | ¢ | III |
| 5／12 | $\mathbf{X}$ | 3 |
| 5／13 | G | 畋 |
| 5／14 | ［1］ | 4 |
| 5／15 | II | H |
| 6／0 | 4 | D |
| 6／10 | j | H |
| 7／1 | Ḱ | я |
| 7／6 | B | $\mathbf{x}$ |
| 7／7 | $\stackrel{r}{1}$ | B |
| 7／8 | b | b |
| 7／9 | や | \％ |
| 7／11 | h | II |
| 7／12 | I | 3 |
| 7／13 | b | III |
| 7／14 | 四 | Y |


| OPTION NUMBER <br> AND <br> C12 C13 C14 | ASSOCIATED ALPHABETS |
| :--- | :--- |
| 1000 | Serbian, Croatian, Macedonian (Cyrillic) |
| see Figure 23 and 25 |  |

See Figure 14 character Set 36 for code in packets with $Y=28$

|  | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 |  | 0 | $t$ | II | 0 | $\pi$ |
| 1 | ! | 1 | A | P | $\alpha$ | p |
| 2 | " | 2 | B | , | $\beta$ | $\leqslant$ |
| 3 | * | 3 | $\Gamma$ | $\Sigma$ | $\gamma$ | $\sigma$ |
| 4 | \$ | 4 | $\Delta$ | T | $\delta$ | $\tau$ |
| 5 | $\%$ | 5 | E | Y | $\epsilon$ | 0 |
| 6 | \& | 6 | 2 | $\Phi$ | $\zeta$ | ¢ |
| 7 | ' | 7 | H | X | $\eta$ | $\chi$ |
| 8 | $($ | 8 | $\theta$ | $\Psi$ | $\theta$ | $\Psi$ |
| 9 | ) | 9 | I | $\Omega$ | $\downarrow$ | $\omega$ |
| 10 | * | : | K | İ | K | i |
| 11 | + | ; | $\wedge$ | $\ddot{\text { Y }}$ | $\lambda$ | ù |
| 12 | , | " | M | ${ }_{\alpha}{ }^{\prime}$ | $\mu$ | 6 |
| 13 | - | $=$ | N | $\dot{\epsilon}$ | $v$ | ט |
| 14 | - | " | $\Xi$ | $\eta$ | $\boldsymbol{\xi}$ | $\dot{\omega}$ |
| 15 | 1 | ? | 0 | $\mathfrak{l}$ | $\bigcirc$ | $\square$ |

Note: The character in position $2 / 0$ is to be interpreted as SPACE


Notes:
Column 4 contains diacritical marks for association with another character addressed to the same presentation position.

The character in position $2 / 0$ is to be interpreted as SPACE.
The character in position $5 / 6$ represents the European Currency symbol.
Blank positions are "transparent", permitting the display of any character otherwise overwritten.

|  | 0 | 1 |  | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | c | D | E | F |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 |  | D | A | ALG | AND | ISR | I | BEL | BLR |  | ALB | AUT | hng | MLT | D | CNR | EgY |
| 2 |  | GRC | c | CYP | SM | SUI | JOR | FNL | Lux | BUL | DNK | GIB | IRQ | G | LBY | Rou | F |
| 3 |  | MR |  | TCH | POL | CVA |  | SYR | TUN | mar | LIE | ISL | MCO |  |  | E | NOR |
| 4 |  |  |  | IRL | TUR |  |  | Yug | UKR | hol |  | LBN |  |  |  | s |  |
| 5 |  |  |  |  |  |  |  |  | URS | POR |  |  |  |  |  |  |  |


| $\begin{aligned} & \text { ALB } \\ & \text { ALG } \end{aligned}$ | Albania <br> Algeria (Peoples | D | Germany (Federal Republic of) | HOL | Netherlands (The Ringdom |
| :---: | :---: | :---: | :---: | :---: | :---: |
| ALG | Democratic | GIB | Gibraltar |  | Of) |
|  | Republic of) | GRC | Greece | NOR | Norway |
| AND | Andorra | HNG | Hungary | POL | Poland |
| AUT | Austria | ISL | Iceland | POR | Portugal |
| AZR | Azores | IRQ | Iraq (Republic | ROU | Roumania |
| BEL | Belgium |  | Of) | SM | San Marino |
| URS | Byelorussia | IRL | Ireland |  | (Republic of) |
| BUL | Bulgaria | ISR | Israel | E | Spain |
| CNR | Canaries | I | Italy | S | Sweden |
| URS | Commonwealth of Independent | JOR | Jordan (Hashemite | SUI | Switzerland (Confederation |
| CYP | Cyprus (Republic of) | LBY | Libya (Social. Peoples Libyan | SYR | Syrian Arab Republic |
| TCH | Czech and Slovak |  | Arab Jamahiriya) | TUN | Tunisia |
|  | Federal Republic | LBN | Lebanon | TUR | Turkey |
| DNK | Denmark | LIE | Liechtenstein | UKR | Ukraine |
| EGY | Egypt (Arab | LUX | Luxembourg | G | United Kingdom |
|  | Republic of) | MDR | Madeira | CVA | Vatican City |
|  | Faroe | MLT | Malta (Republic |  | State |
| FNL <br> $F$ | Finland France | MCO | Of) | YUG | Yugoslavia <br> Socialist |
|  |  | MRC | Morocco (Kingdom of) |  | Federal <br> Republic of) |


| CODE <br> 00 | PROGRAMME TYPE <br> Information not included | REFERENCE KUMBER |
| :---: | :---: | :---: |
| 01-3E | INTENDED AUDIENCE |  |
| 08 | General Audience | 2.0 .0 |
|  | Special Groups |  |
| 10 | Ethnic \& Immigrant Groups | 1.1 .0 |
| 11 | Ethnic Groups | 1.1 .1 |
| 12 | Immigrant Groups | 1.1 .2 |
| 18 | Age Groups | 1.2 .0 |
| 19 | Children (0-13 years) | 1.2.1 |
| 1A | Young people (14 years and over) | 1.2.2 |
| 1 F | Retired people | 1.3 .0 |
| 20 | Disabled people | 1.4 .0 |
| 21 | Blind people | 1.4.1 |
| 22 | Deaf people | 1.4 .2 |
| 28 | Householders | 1.5 .0 |
| 30 | Occupational Status Groups | 1.6 .0 |
| 31 | Unemployed people | 1.6 .1 |
| 32 | Studenta | 1.6 .2 |
| 33 | Farmers | 1.6 .3 |
| 34 | Fishermen \& Sailors | 1.6 .4 |
| 38 | Travellers | 1.7 .0 |
| 39 | Motorists | 1.7 .1 |
| 3A | Tourists | 1.7 .2 |
| 3F | ALARM/EMERGENCY IDENTIFICATION |  |
| 40-7F | CONTENT |  |
| 40 | Public affairs |  |
| 41 | General Domestic | 1.1 .0 |
| 42 | Legal and social | 1.2 .0 |
| 43 | Economic, industrial \& financial | 1.3 .0 |
| 44 | Housing, environment \& health | 1.4 .0 |
| 45 | Communication | 1.5 .0 |
| 46 | Educational \& cultural | 1.6 .0 |
| 47 | International relations \& defence | 1.7 .0 |
| 48 | Science \& the humanities | 2.0 .0 |
| 49 | Natural sciences | 2.1 .0 |
| 4A | Social sciences | 2.2 .0 |
| 4B | Humanities | 2.3 .0 |
| 4 C | Other sciences or humanities | 2.9 .0 |

continued in part 2

| 50 | Music | 3.1 .0 |
| :---: | :---: | :---: |
| 51 | Serious | 3.1.1 |
| 52 | Light classical | 3.1.2 |
| 53 | Light | 3.1 .3 |
| 54 | Jazz | 3.1.4 |
| 55 | Folk | 3.1 .5 |
| 56 | Rock |  |
| 57 | Other music | 3.1 .9 |
| 58 | Drama, arts | 3.0 .0 |
| 5A | Ballet \& dance | 3.2 .0 |
| 5B | Drama | 3.3.0 |
| 5C | Literature/poetry | 3.4.0 |
| 5D | Media affairs | 3.5 .0 |
| 5E | Painting, sculpture, architecture | 3.6 .0 |
| 5 F | Other drama, arts | 3.9 .0 |
| 60 | Philosophies of life | 4.0 .0 |
| 61. | Christian religion | 4.1 .0 |
| 62 | Non-Christian religion | 4.2 .0 |
| 63 | Non-religious philosophies of life | 4.3.0 |
| 67 | Other philosophies of life | 4.9 .0 |
| 68 | Sports | 5.0.0 |
| 69 | Non-instrumental ball games | 5.1 .0 |
| 6A | Instrumental ball games | 5.2.0 |
| 6B | Winter sports | 5.3.0 |
| 6 C | Water sports | 5.4.0 |
| 6D | Racing and equestrian sports | 5.5.0 |
| 6 E | Athletics | 5.6.0 |
| 6F | Martial arts | 5.7 .0 |
| 70 | Leisure \& hobbies | 6.0 .0 |
| 71 | Do-it-yourself | 6.1 .0 |
| 72 | Gardening | 6.2 .0 |
| 73 | Tourism | 6.3 .0 |
| 74 | Keep fit | 6.4 .0 |
| 77 | other leisure or hobbies | 6.9 .0 |
| 78 | Light entertainment, folklore and human interest | 7.0 .0 |
| 7A | Light entertainment | 7.1 .0 |
| 7 B | Folklore/festivities | 7.2 .0 |
| 7 C | Human interest | 7.2 .0 |
| 7 F | Other light entertainment | 7.9 .0 |
| $80-\mathrm{FF}$ | SERVICE SPECIFIC SERIES CODES |  |

This code table is based on the EBU classification system 'ESCORT' [13].



CODE TABLE FOR ENCRYPTION METHODS

| Method <br> Number | Bits 1 to 8 <br> in Transmission Order | Name of Encryption Method |
| :---: | :---: | :--- |
| 1 | 00000000 | Variable Length Algorithm <br> Using a One Way Function <br> (see Figures 35 and 36) |
| 2 | 10000000 | Block Encipherment Algorithm <br> Using Differential Code Book <br> or Output Feedback <br> (see note below and Figure 39) <br> Method Not Specified |
| Figure 33 (a) REGISTER OF ENCRYPTION METHODS |  |  |

CODE TABLE FOR SCRAMBLING METHODS

| Method <br> Number | Bits l to 5 <br> in Transmission Order | Name of Scrambling Method |
| :---: | :---: | :--- |
| 1 | 00000 | Variable Length Algorithm <br> Using a One Way Function <br> Method A <br> (see Figurea 35 and 36) <br> Block Encipherment Algorithm <br> Using Output Feedback |
| 3 | 10000 | Variable Length Algorithm <br> Using a One Way Function <br> Method B |
| 4 | 11000 | MAC Scrambler <br> Method Not Specified |
| Figure 33 (b) REGISTER OF SCRAMBLING METHODS |  |  |

Note: Differential Code Book The first 8 bytes of conversion use the ISO Electronic Code Book [10]. For subsequent 8-byte blocks, the input data at decipherment is first exclusive-oRed' with the result of the previous conversion and then converted using the Electronic Code Book.

Figures 33 (a) and (b) REGISTER OF ENCRYPTION AND SCRAMBLING METBODS




NOT ENCRYPTED

msb

Figure 35
VARIABLE LENGTE ENCRYPTION ALGORITHM



| e7ed дesn patquresos | 87798 Jequinn aวuənbas | 87798 <br> дaqunn uotzestftzuapi <br> องไฺภอง | $\begin{array}{r} 87798 \\ 20 / 70 \\ \text { edKJ צコOTg } \end{array}$ | $\begin{array}{r} 87798 \\ 470 \end{array}$ |
| :---: | :---: | :---: | :---: | :---: |

## 

（q）$\llcorner\varepsilon$ exnbta



|  | e780 ג380 | абеssak pa7dイx | pazdरxכug zon abessaw |  | 97a |
| :---: | :---: | :---: | :---: | :---: | :---: |


| DLE <br> 8 bits | Block Type 04/03 <br> 8 bits | Encryption Method 8 bits | New system Key Label 8 bits | Current system Key Label <br> 8 bits | New System Key <br> 56 bits | Current System Rey 56 bits |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| DATA NOT ENCRYPTED |  |  |  |  | ENCRYPTED W | NEW SYSTEM KEY |
| Figure $37(\mathrm{~d})$ SYSTEM KEY BLOCK MESSAGE: (See Section 24.3.4) |  |  |  |  |  |  |


| DLE <br> 8 bits | Block Type 04/04 <br> 8 bits | Shared User Address 20 bits | set <br> to 0 4 bits | New System Key 56 bits | User Enabling Bits <br> 152 bits |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| DATA NOT ENCRYPTED |  |  |  | ENCRYPTED WITH SHARED DISTRIBUTION KEY |  |  |
| Figure $37(\mathrm{e})$ |  |  | SHARED USER MESSAGE BLOCK: (See Section 24.3.5) |  |  |  |


| DLE <br> 8 bits | ```Block Type 04/05 8 bits``` | Unique User Address 32 bits | Service <br> Mode <br> 2 bits | ```Service Reference Number 2 bits``` |
| :---: | :---: | :---: | :---: | :---: |
| NOT ENCRYPTED |  |  | ENCRYPTED WITH UNIQUE KEY |  |


| Current or New Shared Address 20 bits | Unique <br> Equipment Key 56 bits | ```7 Service Numbers 8 Bits Each 56 bits``` | Current or New Shared Distribution Rey 56 bits | Current or New User Enabling Bit Position 8 bits |
| :---: | :---: | :---: | :---: | :---: |
| ENCRYPTED WITH UNIQUE KEY |  |  |  |  |

Figure $37(f)$ UNIQUE USER BLOCK: (See Section 24.3.6)



Figure 37(h) SERVICE ADDRESS BLOCK: PAGE FORMAT SERVICES (See Section 24.3.8)




Figure $40(\mathrm{~h})$ UNIQUE SYSTEM KEY MESSAGE BLOCK (See Section 24.4.9)


Figure 40(j) Shared equipment rex message block (See section 24.4.10)


Figure $40(k)$ OVER-AIR CREDIT MESSAGE BLOCK (See Section 24.4.11)




[^0]:    NOTE In all cases the LEAST SIGNIFICANT BIT is transmitted first.
    $P=$ Hamming $(8,4)$ Protection bit, $D=$ Data bit in sequence PDPDPDPD in any byte.
    Odd parity bytes are coded DDDDDDDP in any byte

